# Learning Go <br>  

Thanks to:<br>Go Authors Google<br>Go Nuts mailing list

## (i) (ㅇ) (ㅇ)

Learning as we Go.
Updated to Go release.2010-10-20.

## Table of Contents

1 Introduction ..... vi
Official documentation ..... vii
Getting Go ..... vii
Origins ..... viii
Exercises ..... ix
Answers ..... xi
2 Basics ..... 1
Hello World ..... 2
Compiling and running code ..... 2
Variables, types and keywords ..... 3
Operators and built-in functions ..... 6
Go keywords ..... 7
Control structures ..... 8
Arrays, slices and maps ..... 13
Exercises ..... 15
Answers ..... 17
3 Functions ..... 20
Scope ..... 21
Multiple return values ..... 22
Named result parameters ..... 23
Deferred code ..... 24
Variadic parameters ..... 25
Functions as values ..... 26
Panic and recovering ..... 26
Exercises ..... 26
Answers ..... 29
4 Packages ..... 34
Building a package ..... 35
Identifiers ..... 36
Initialization ..... 37
Documenting packages ..... 37
Testing packages ..... 38
Useful packages ..... 39
Exercises ..... 41
Answers ..... 43
5 Beyond the basics ..... 46
Allocation ..... 47
Defining your own ..... 49
Conversions ..... 50
Exercises ..... 52
Answers ..... 55
6 Interfaces ..... 58
Interface names ..... 60
Introspection ..... 60
Exercises ..... 64
Answers ..... 65
7 Concurrency ..... 66
Exercises ..... 68
Answers ..... 71
8 Communication ..... 74
Files ..... 74
Command line arguments ..... 75
Executing commands ..... 75
Exercises ..... 76
Answers ..... 79
A Exercises ..... 82
Exercises ..... 82
Answers ..... 85
B Colophon ..... 90
Contributors ..... 90
Further reading ..... 90
C Index ..... 92
D Bibliography ..... 94
List of Figures
1.1 Chronology of Go ..... ix
2.1 Array versus slice ..... 14
3.1 A simple LIFO stack ..... 27
5.1 Pointers and types ..... 46
6.1 Pealing away the layers using reflection ..... 63
List of Tables
2.1 Operator precedence ..... 6
2.2 Pre-defined functions in Go ..... 6
2.3 Keywords in Go ..... 7
5.1 Valid conversions ..... 51

## List of Code Examples

2.1 Hello world ..... 2
2.2 Makefile for a program ..... 3
2.3 Declaration with = ..... 3
2.4 Declaration with ..... 3
2.5 Familiar types are still distinct ..... 4
2.6 Arrays and slices ..... 14
2.7 Simple for-loop ..... 17
2.8 Loop calls function ..... 17
2.9 Fizz-Buzz ..... 17
2.10 Strings ..... 18
2.11 Runes in strings ..... 19
2.12 Reverse a string ..... 19
3.1 A function declaration ..... 20
3.2 Local scope ..... 21
3.3 Global scope ..... 21
3.4 Scope when calling functions from functions ..... 21
3.5 Without defer ..... 24
3.6 With defer ..... 24
3.7 Function literal ..... 25
3.8 Access return values within defer ..... 25
3.9 Variadac function declaration ..... 25
3.10 Anonymous function ..... 26
3.11 Functions as values in maps ..... 26
3.12 The push and pop functions ..... 30
3.13 Stack usage ..... 30
3.14 stack.String() ..... 30
3.15 A function with variable number of arguments ..... 31
3.16 Fibonacci function in Go ..... 31
3.17 A Map function ..... 32
4.1 A small package ..... 34
4.2 Use of the even package ..... 34
4.3 Makefile for a package ..... 35
4.4 Test file for even package ..... 38
4.5 Stack in a package ..... 43
4.6 A (rpn) calculator ..... 43
5.1 Make use of a pointer ..... 46
5.2 Dereferencing a pointer ..... 47
5.3 Structures ..... 49
5.4 A generic map function in Go ..... 55
5.5 A cat program ..... 56
6.1 Defining a struct and methods on it ..... 58
6.2 A function with a empty interface argument ..... 59
6.3 Failing to implement an interface ..... 59
6.4 Failure extending built-in types ..... 60
6.5 Failure extending non-local types ..... 60
6.6 Dynamically find out the type ..... 60
6.7 A more generic type switch ..... 61
6.8 Introspection using reflection ..... 62
6.9 Reflection and the type and value ..... 62
6.10 Reflect with private member ..... 64
6.11 Reflect with public member ..... 64
7.1 Go routines in action ..... 66
7.2 Go routines and a channel ..... 67
7.3 Using select ..... 68
7.4 Channels in Go ..... 71
7.5 Adding an extra quit channel ..... 71
7.6 A Fibonacci function in Go ..... 72
8.1 Reading from a file (unbufferd) ..... 74
8.2 Reading from a file (bufferd) ..... 74
8.3 Processes in Perl ..... 76
8.4 Processes in Go ..... 79
A. 4 uniq(1) in Perl ..... 83
A. $1 \mathrm{wc}(1)$ in Go ..... 85
A. 2 Average function in Go ..... 86
A. 3 Number cruncher ..... 86
A. 5 uniq(1) in Go ..... 89
List of Exercises
1 (1) Documentation ..... ix
2 (1) For-loop ..... 15
3 (1) FizzBuzz ..... 15
4 (1) Strings ..... 15
5 (3) Integer ordering ..... 26
6 (4) Scope ..... 26
7 (5) Stack ..... 27
8 (5) Var args ..... 27
9 (5) Fibonacci ..... 27
10 (4) Map function ..... 27
11 (2) Stack as package ..... 41
12 (7) Calculator ..... 41
13 (6) Map function with interfaces ..... 52
14 (6) Pointers ..... 52
15 (5) Pointers and reflection ..... 52
16 (6) Linked List ..... 52
17 (6) Cat ..... 52
18 (8) Method calls ..... 53
19 (6) Interfaces and compilation ..... 64
20 (4) Channels ..... 68
21 (7) Fibonacci II ..... 68
22 (8) Processes ..... 76
23 (3) Minimum and maximum ..... 82
24 (5) Bubble sort ..... 82
25 (5) Word and letter count ..... 82
26 (4) Average ..... 82
27 (9) Number cruncher ..... 82
28 (4) Uniq ..... 83

I am interested in this and hope to do something.

On adding complex numbers to Go KEN THOMPSON

This is an introduction to the Go language from Google. Its aim is to provide a guide to this new and innovative language. What is Go? From the website [6]:

> The Go programming language is an open source project to make programmers more productive. Go is expressive, concise, clean, and efficient. Its concurrency mechanisms make it easy to write programs that get the most out of multicore and networked machines, while its novel type system enables flexible and modular program construction. Go compiles quickly to machine code yet has the convenience of garbage collection and the power of run-time reflection. It's a fast, statically typed, compiled language that feels like a dynamically typed, interpreted language.

The intended audience of this book is people who are familiar with programming and know multiple programming languages, be it C[14], $\mathrm{C}^{++}$[22], Perl [16], Java [15], Erlang[13], Scala[1] or Haskell[7]. This is not a book which teaches you how to program, this is a book that just teaches you how to use Go.

As with learning new things, probably the best way to do this is to discover it for yourself by creating your own programs. Therefor includes each chapter a number of exercises (and answers) to acquaint you with the language. An exercise is numbered as $\mathbf{Q} n$, where $n$ is a number. After the exercise number another number in parentheses displays the difficulty of this particular assignment. This difficulty ranges from 0 to 2 , where 0 is easy and 2 is difficult. Then a short name is given, for easier reference. For example

## Q1. (1) A map function ...

introduces a question numbered $\mathbf{Q 1}$ of a level 1 difficulty, concerning a map ()-function. The answers are included after the exercises on a new page, except for those pesky unresolved research problems. The numbering and setup of the answers is identical to the exercises, except that an answer starts with $\mathbf{A} n$, where the number $n$ corresponds with the number of the exercise.

Go is a young language, where features are still being added or even removed. It may be possible that some text is outdated when you read it. Some exercise answers may become incorrect as Go continues to evolve. We will do our best to keep this document up to date with respect to the latest Go release. An effort has been made to create "future proof" code examples.

The following convention is used throughout this book:

- Code is displayed in DejaVu Mono;
- Keywords are displayed in DejaVu Mono Bold;
- Comments are displayed in DejaVu Mono Italic;
- Extra remarks in the code $\leftarrow$ Are displayed like this;
- Longer remarks get a number - 1 - with the explanation following;
- Line numbers are printed on the right side;
- Shell examples use a \% as prompt;
- An emphasized paragraph is indented and has a vertical bar on the left.


## Official documentation

There already is a substantial amount of documentation written about Go. The Go Tutorial [5], and the Effective Go document [2]. The website http://golang.org/doc/ is a very good starting point for reading up on $\mathrm{Go}^{a}$. Reading these documents is certainly not required, but is recommended.

Go comes with its own documentation in the form of a Go program called godoc. You can use it yourself to look in the online documentation. For instance, suppose we want to know more about the package hash. We would then give the command godoc hash. How to read this and how you can create your own package documentation is explained in chapter 4.

## Getting Go

There are currently (2010) no packages for Go in any Linux distribution. The route to install Go is thus slightly longer than it could be. When Go stabilizes this situation will change probably. For now, you need to retrieve the code from the mercurial archive and compile Go yourself. For other Unix like systems the procedure is the same

- First install Mercurial (to get the hg command). In Ubuntu/Debian/Fedora you must install the mercurial package;
- For building Go you also need the packages: bison, gcc, libc6-dev, ed, gawk and make;
- Set the environment variable GOROOT to the root of your Go install:
\% export GOROOT=~/go
- Then retrieve the Go source code:
\% hg clone -r release https://go.googlecode.com/hg/ \$GOROOT
- Set your PATH to so that the Shell can find the Go binaries: \% export PATH=\$GOROOT/bin:\$PATH
- Compile Go
\% cd \$GOROOT/src
\% ./all.bash

If all goes well, you should see the following:

[^0]TODO
We don't explain how to read this in that chapter.

Installed Go for linux/amd64 in /home/gobook/go.
Installed commands in /home/gobook/go/bin.
The compiler is 6 g .

You now have Go installed on your system and you can start playing.

## Keeping up to date

New releases are announced on the Go Nuts mailing list [12]. To update an existing tree to the latest release, you can run:

```
% cd $GOROOT
% hg pull
% hg update release
% cd src
% ./all.bash
```

To see what you are running right now:
\% cd \$GOROOT
\% hg identify
79997f0e5823 release/release.2010-10-20

That would be release.2010-10-20.

## Origins

Go has it origins in Plan 9 [10]. Plan 9 is (or was supposed to be) the successor of Unix. As you may know one of the core ideas of Unix is "everything is a file", so the read () and write() system calls work as well on normal files as they do on I/O devices. However for some devices this has never happened, notably network and video devices. In Plan 9 this "everything is a file"-idea is taken to the next level and truly everything is presented to the user as a file, including networking and video devices. This has nothing to do with Go per se, but Plan 9 included a language called Limbo [9]. Quoting from the Limbo paper:

Limbo is a programming language intended for applications running distributed systems on small computers. It supports modular programming, strong type checking at compile- and run-time, inter process communication over typed channels, automatic garbage collection, and simple abstract data types. It is designed for safe execution even on small machines without hardware memory protection.

That sounds a lot like Go, and one could say Go is a reimplementation of Limbo on Unix-like systems like Linux. One feature of Limbo that is also included in Go is the support for cross compiling.

Another feature Go inherited from Limbo is channels (see chapter 7). Again from the Limbo documentation.
[A channel] is a communication mechanism capable of sending and receiving objects of the specified type to another agent in the system. Channels may be used to communicate between local processes; using library procedures, they may be connected to named destinations. In either case send and receive operations may be directed to them.

The channels in Go are easier to use than those in Limbo. If we dig even deeper in the history of Go we also find references to "Newsqueak" [21], which pioneered the use of channel communication in a C-like language. Channel communication isn't unique to these languages, a big non-C-like language which also uses them is Erlang [13].

Figure 1.1. Chronology of Go


The whole of idea of using channels to communicate with other processes is called Communicating Sequential Processes (CSP) and was conceived by C. A. R. Hoare [17], who incidentally is the same man that invented QuickSort [18].

Go is the first C-like language that is widely available, runs on many different platforms and makes concurrency easy (or easier).


## Exercises

Q1. (1) Documentation

1. Go's documentation can be read with the godoc program, which is included the Go distribution.
godoc hash gives information about the hash package. Reading the documentation on container gives the following result:

SUBDIRECTORIES

## heap

list
ring
vector

With which godoc command can you read the documentation of vector contained in container?

## Answers

A1. (1) Documentation

1. The package vector is in a subdirectory of container, so you will only need godoc container/vector.
Specific functions inside the "Go manual" can also be accessed. For instance the function Printf is described in fmt, but to only view the documentation concerning this function use: godoc fmt Printf.
You can even display the source code with: godoc -src fmt Printf.

## Basics

In Go, the code does exactly what it says on the page.

Go Nuts mailing list
ANDREW GERRAND
There are a few things that make Go different from most other languages out there.

## Clean and Simple

Go strives to keep things small and beautiful, you should be able to do a lot in only a few lines of code;

## Concurrent

Go makes it easy to "fire off" functions to be run as very lightweight threads. These
threads are called goroutines ${ }^{a}$ in Go;

## Channels

Communication with these goroutines is done via channels [26][17];
Fast
Compilation is fast and execution is fast. The aim is to be as fast as C . Compilation time is measured in seconds;

## Safe

Go has garbage collection, no more free( ) in Go, the language takes care of this;

## Standard format

A Go program can be formatted in (almost) any way the programmers want, but an official format exist. The rule is very simple: The output of the filter gofmt is the official endorsed format.

## Postfix types

Types are given after the variable name, thus var a int, instead of int a; as one would in C;

UTF-8
UTF-8 is everywhere, in strings and in the program code. Finally you can use $\Phi=$ $\Phi+1$ in your source code;

## Open Source

The Go license is completely open source, see the file LICENSE in the Go source code distribution;

Fun
Programming with Go should be fun!
Erlang [13] also shares some of the features of Go. Notable differences between Erlang and Go is that Erlang borders on being a functional language, where Go is an imperative one. And Erlang runs in a virtual machine, while Go is compiled. Go also has a much more Unix-like feeling to it.

[^1]
## Hello World

In the Go tutorial，Go is presented to the world in the typical manner：letting it print＂Hello World＂（Ken Thompson and Dennis Ritchie started this when they presented the C language in the nineteen seventies）．We don＇t think we can do better，so here it is，＂Hello World＂in Go．

## Listing 2．1．Hello world

```
package main 11
import "fmt" // Implements formatted I/O. 2 3
/* Print something */ 3 5
func main() {4 6
5 7
    fmt.Printf("Hello, world; or \kappa\alpha\lambda\eta\mu\epsiloń\rho\alpha \kappaó\sigma\mu\epsilon; or こんにちは世界\n") 8
}

Lets look at the program line by line．
（1）This first line is just required．All Go files start with package＜something＞，package main is required for a standalone executable；
（2）This says we need＂fmt＂in addition to main．A package other than main is commonly called a library，a familiar concept of many programming languages（see chapter 4）． The line ends with a comment which is started with／／；
（3）This is also a comment，but this one is enclosed in／＊and＊／；
（4）Just as package main was required to be first，import may come next．In Go，package is always first，then import，then everything else．When your Go program is exe－ cuted，the first function called will be main．main（），which mimics the behavior from C．Here we declare that function；
（5）On line 8 we call a function from the package \(f m t\) to print a string to the screen．The string is enclosed with＂and may contain non－ASCII characters．Here we use Greek and Japanese．

\section*{Compiling and running code}

The Go compiler is named＜number＞g，where the number is 6 for 64 bit Intel and 8 for 32 bit Intel．The linker has a similar naming scheme：＜number＞l．In this book we will use 6 g and 61 for all the compiles．To compile the code from above，we use：
\％6g helloworld．go

And then we link it with 61 ：
\％ \(6 l\) helloworld． 6

And then we run it：
```

%./6.out }\leftarrow\mathrm{ The default name for a (64 bit) Go executable
Hello, world; or \kappa\alpha\lambda\eta\mu\epsiloń\rho\alpha \kappaó\sigma\mu\epsilon; or こんにちは世界

```

\section*{Using a Makefile}

Another，less laborious（once setup），way to build a Go program，is to use a Makefile．The following one can be used to build helloworld：

\section*{Listing 2．2．Makefile for a program}
```


# Copyright 2009 The Go Authors. All rights reserved. 1

# Use of this source code is governed by a BSD-style 2

# license that can be found in the LICENSE file. 3

include \$(GOROOT)/src/Make.inc 5
TARG=helloworld 7
GOFILES=\ 8
helloworld.go\ 9
include \$(GOROOT)/src/Make.cmd 11

```

At line 7 you specify the name for your compiled program and on line 9 you enumerate the source files．Now only an invocation of make is enough to get your program compiled．Note that Go ships with the variant of make，called gomake，which is（currently）a small wrapper around make．As the build system for Go programs may change in the future and make make go away，we use gomake．Note that this Makefile creates an executable named helloworld， not 6 ．out．

\section*{Variables，types and keywords}

In the next sections we will look at variables，basic types，keywords and control structures of our new language．

Go is different from other languages in that the type of a variable is specified after the variable name．So not：int a，but a int．When declaring a variable it is assigned the ＂natural＂null value for the type．This means that after var a int，a has a value of 0 ．With var s string，s is assigned the zero string，which is＂＂．

Declaring and assigning in Go is a two step process，but they may be combined．Com－ pare the following pieces of code which have the same effect．

Listing 2．3．Declaration with \(=\)
Listing 2．4．Declaration with ：＝
var a int
a ：＝ 15
var \(b\) bool
b ：＝false
a \(=15\)
b＝false

On the left we use the var keyword to declare a variable and then assign a value to it．The code on the right uses ：＝to do this in one step（this form may only be used inside functions）．In that case the variable type is deduced from the value．A value of 15 indicates
an int, a value of false tells Go that the type should be bool. Multiple var declarations may also be grouped, const and import also allow this. Note the use of parentheses:
```

var (
x int
b bool
)

```

Multiple variables of the same type can also be declared on a single line: var \(x, y\) int, makes \(x\) and \(y\) both int variables. You can also make use of parallel assignment:
a, b := 20, 16
Which makes a and b both integer variables and assigns 20 to a and 16 to b .
A special name for a variable is _(underscore). Any value assigned to it, is discarded. In this example we only assign the integer value of 35 to b and discard the value 34 .
```

_, b := 34, 35

```

Declared, but otherwise unused variables are a compiler error in Go, the following code generates this error: i declared and not used
```

package main
func main() {
var i int
}

```

\section*{Boolean types}

A Boolean type represents the set of Boolean truth values denoted by the predeclared constants true and false. The Boolean type is bool.

Numerical types

Go has the well known types such as int and float. These types have the appropriate length for your machine, meaning that on a 32 bits machine they are 32 bits, and on a 64 bits machine they are 64 bits.

If you want to be explicit about the length you can have that too with int32, or uint32. The full list for (signed and unsigned) integers is int8, int16, int32, int64 and byte, uint8, uint16, uint32, uint64. With byte being an alias for uint8. For floating point values there only is float 32 and float 64 .

Note however that these types are all distinct and assigning variables which mix these types is a compiler error, like in the following code:

\section*{Listing 2.5. Familiar types are still distinct}


Gives the error on the assignment on line 7:
types.go:7: cannot use a + a (type int) as type int32 in assignment
When assigning values octal, hexadecimal and the scientific notation may be used: 077, \(0 x F F, 1 e 3\) or \(6.022 e 23\) are all valid.

\section*{Constants}

Constants in Go are just that - constant. They are created at compile time, and can only be numbers, strings or booleans; const \(x=42\) makes \(x\) a constant. You can use iota \({ }^{\boldsymbol{b}}\) to enumerate values.
```

const (
a = iota
b = iota
)

```

The first use of iota will yield 0 , so a is equal to 0 , whenever iota is used again on a new line its value is incremented with 1 , so \(b\) has a value of 1 .

You can even do the following, let Go repeat the use of iota:
```

const (
a = iota
b // implicitly 'b = iota'
)

```

Strings

An important other built in type is string. Assigning a string is as simple as:
```

s := "Hello World!"

```

Strings in Go are a sequence of UTF-8 characters enclosed in double quotes ("). If you use the single quote (') you mean one character (encoded in UTF-8) - which is not a string in Go.

Once assigned to a variable the string can not be changed anymore: strings in Go are immutable. For people coming from C, the following is not legal in Go:
```

var s string = "hello"
s[0] = 'c' \leftarrow Change first char. to 'c', this is an error

```

To do this in Go you will need the following:
```

s := "hello"
c := []byte(s)
c[0] = 'c'
s2 := string(c)
fmt.Printf("%s\n", s2)

```
(1) Convert s to an array of bytes, see chapter 5 section "Conversions";

\footnotetext{
\({ }^{\boldsymbol{b}}\) The word [iota] is used in a common English phrase, 'not one iota', meaning 'not the slightest difference', in reference to a phrase in the New Testament: "until heaven and earth pass away, not an iota, not a dot, will pass from the Law." [28]
}
(2) Change the first element of this array;
(3) Create a new string s2 with the alteration;
(4) print the string with fmt. Printf.

\section*{Complex numbers}

Go has native support for complex numbers. If you use them you need a variable of the type complex. If you need to specify the number of bits you have complex32 and complex64 for 32 and 64 bits. Complex numbers are written as re \(+i m i\), where re is the real part, im is the imaginary part and \(i\) is the literal ' \(i\) ' \((\sqrt{-1})\). An example of using complex numbers:
var c complex \(=5+5 i\); fmt. Printf("Value is: \%v", c")
will print: ( \(5+5 i\) )

\section*{Operators and built-in functions}

Go supports the normal set of numerical operations, table 2.1 lists the current ones and their relative precedence. They all associate from left to right.

Table 2.1. Operator precedence
\begin{tabular}{ll} 
Precedence & Operator(s) \\
\hline Highest & \(* / \% \ll \gg \& \&^{\wedge}\) \\
& \(+-\mid \wedge\) \\
& \(==!=\ll \ggg\) \\
& \(<-\) \\
& \(\& \&\) \\
& \\
& \\
& \\
&
\end{tabular}

\begin{abstract}
\(+-* /\) and \(\%\) all do what you would expect, \& | ^ and \(\&^{\wedge}\) are bit operators for bit-wise and, or, xor and bit clear respectively. Although Go does not support operator overloading (or method overloading for that matter), some of the built-in operators are overloaded. For instance + can be used for integers, floats, complex numbers and strings (adding strings is concatenating them).

A small number of functions are predefined, meaning you don't have to include any package to get access to them. Table 2.2 lists them all.
\end{abstract}

Table 2.2. Pre-defined functions in Go
\begin{tabular}{llll} 
close & new & panic & cmplx \\
closed & make & panicnl & real \\
len & copy & print & imag \\
cap & & printnl &
\end{tabular}
close and closed are used in channel communication and the closing of those channels, see chapter 7 for more on this.
section needs to be written
len and cap are used on a number of different types, len is used for returning the length of strings and the length of slices and arrays. See section "Arrays, slices and maps" for the details of slices and arrays and the function cap.
new is used for allocating memory for user defined data types. See section "Allocation with new" on page 47.
make is used for allocating memory for built-in types (maps, slices and channels). See section "Allocation with make" on page 47.
copy is used for copying slices. new, make and copy make their appearance in chapter 5.
panic and panicln are used for an exception mechanism. See the section "Panic and recovering" on page 26 for more.
print and println are low level printing functions that can be used without reverting to the fmt package. These are mainly used for debugging.
cmplx, real and imag all deal with complex numbers.

\section*{Go keywords}

TODO
copy isn't handled
there
TODO
Needs to be written

\section*{TODO}

Refer to other documentation.

\section*{Table 2.3. Keywords in Go}
\begin{tabular}{lllll} 
break & default & func & interface & select \\
case & defer & go & map & struct \\
chan & else & goto & package & switch \\
const & fallthrough & if & range & type \\
continue & for & import & return & var
\end{tabular}

Table 2.3 lists all the keywords in Go. In the following paragraphs and chapters we will cover them. Some of these we have seen already.
- For var and const see section "Variables, types and keywords" on page 3;
- package and import are briefly touch upon in section " ". in chapter 4 they are documented in more detail.

Others deserve more text and have their own chapter/section:
- func is used to declare functions;
- return is used to return from functions, for both func and return see chapter 3 for the details;
- go is used for concurrency (chapter 7);
- interface, see XXX;
- select, see XXX;
- struct, see XXX;
- type, see XXX.

Syntax

Go has a C-like feel when it comes to its syntax. If you want to put two (or more) statements on one line, they must be separated with a semicolon (';'), like so:
```

a = 5; a = a + 1

```

\section*{Control structures}

There are only a few control structures in Go \({ }^{c}\). For instance there is no do or while loop, only a for. There is a (flexible) switch statement and if and switch accept an optional initialization statement like that of for. There also is something called a type switch and a multiway communications multiplexer, select (see chapter 7). The syntax is different (from that in C): parentheses are not required and the bodies must always be bracedelimited.

If

In Go an if looks like this:
```

if x > 0 { }\leftarrow{\mathrm{ is mandatory
return y
} else {
return x
}

```

Mandatory braces encourage writing simple if statements on multiple lines. It is good style to do so anyway, especially when the body contains a control statement such as a return or break.

Since if and switch accept an initialization statement, it's common to see one used to set up a (local) variable.
```

if err := file.Chmod(0664); err != nil { \leftarrow nil is like C's NULL
log.Stderr(err) \leftarrow Scope of err is limited to if's body
return err
}

```

In the Go libraries, you will find that when an if statement doesn't flow into the next statement-that is, the body ends in break, continue, goto, or return, the unnecessary else is omitted.
```

f, err := os.Open(name, os.0_RDONLY, 0)
if err != nil {
return err
}
doSomething(f)

```

\footnotetext{
\({ }^{c}\) This section is copied from [2]
}

This is a example of a common situation where code must analyze a sequence of error possibilities. The code reads well if the successful flow of control runs down the page, eliminating error cases as they arise. Since error cases tend to end in return statements, the resulting code needs no else statements.
```

f, err := os.Open(name, os.O_RDONLY, 0)
if err != nil {
return err
}
d, err := f.Stat()
if err != nil {
return err
}
doSomething(f, d)

```

Syntax wise the following is illegal in Go:
```

if err != nil
{ }\leftarrow\mathrm{ Must be on the same line as the if
return err
}

```

See [2] section "Semicolons" for the deeper reasons behind this.

Goto

Go has a goto statement - use it wisely. With goto you jump to a label which must be defined within the current function. The labels are case sensitive. For instance a loop in disguise:
```

func myfunc() {
i := 0
HERE: }\quad\leftarrow\mathrm{ First word on a line ending with a colon is a label
println(i)
i++
goto HERE }\leftarrow\mathrm{ Jump
}

```

\section*{For}

The Go for loop has three forms, only one of which has semicolons.
```

// Like a C for
for init; condition; post { }
// Like a while
for condition { }
// Like a C for(;;) (endless loop)
for { }

```

Short declarations make it easy to declare the index variable right in the loop.
```

sum := 0
for i := 0; i < 10; i++ \{
sum += i $\leftarrow$ Short for sum $=$ sum $+i$
$\} \quad \leftarrow i$ ceases to exist after the loop

```

Finally, since Go has no comma operator and ++ and -- are statements not expressions, if you want to run multiple variables in a for you should use parallel assignment.
```

// Reverse a
for i, j := 0, len(a)-1; i < j; i, j = i+1, j-1 { }\leftarrow\mathrm{ Parallel assignment
a[i], a[j] = a[j], a[i] \leftarrow Here too
}

```

Break and continue

With break you can quit loops early, break breaks the current loop.
```

for i := 0; i < 10; i++ {
if i > 5 {
break}\leftarrow\mathrm{ Stop this loop, making it only print 0 to 5
}
println(i)
}

```

With loops within loops you can specify a label after break. Making the label identify which loop to stop:
```

J: for j := 0; j < 5; j++ {
for i := 0; i < 10; i++ {
if i > 5 {
break J }\leftarrow\mathrm{ Now it breaks the j-loop, not the i one
}
println(i)
}
}

```

With continue you begin the next iteration of the loop, skipping any remaining code. In the same way as break, continue also accepts a label. The following prints 0 to 5 .
```

for i := 0; i < 10; i++ {
if i > 5 {
continue }\leftarrow\mathrm{ Skip any remaining code
}
println(i)
}

```

Range

The keyword range can be used for loops. It can loop over slices, arrays, string, maps and channels (see chapter 7, section "??"). range is an iterator, that when called, returns a keyvalue pair from the thing it loops over. Depending on what that is, range returns different things.

When looping over a slice or array range returns the index in the slice as the key and value belonging to that index. Consider this code:
```

list := []string{"a", "b", "c", "d", "e", "f"}
for k, v := range list { (2)
// do what you want with }k\mathrm{ and v
}

```
(1) Create a slice (see 2 on page 13) of strings.
(2) Use range to loop over them. With each iteration range will return the index as int and the key as a string, starting with 0 and "a".

You can also use range on strings directly. Then it will break out the individual Unicode characters by parsing the UTF-8 (erroneous encodings consume one byte and produce the replacement rune \({ }^{d}\) U+FFFD). The loop:
```

for pos, char := range "a\Phix" {
fmt.Printf("character '%c' starts at byte position %d\n", char, pos)
}

```
prints
```

character 'a' starts at byte position 0
character '\Phi' starts at byte position 1
character 'x' starts at byte position 3

```

\section*{Switch}

Go's switch is very flexible. The expressions need not be constants or even integers, the cases are evaluated top to bottom until a match is found, and if the switch has no expression it switches on true. It's therefore possible - and idiomatic - to write an if-else-ifelse chain as a switch.
```

func unhex(c byte) byte {
switch {
case '0'<= c \&\& c <= '9':
return c - '0'
case 'a' <= c \&\& c <= 'f':
return c - 'a' + 10
case 'A' <= C \&\& C <= 'F':
return C - 'A' + 10
}
return 0
}

```

There is no automatic fall through, you can however use fallthrough to make do just that. Without fallthrough:

\footnotetext{
\({ }^{\boldsymbol{d}}\) In the UTF-8 world characters are sometimes called runes. Mostly, when people talk about characters, they mean 8 bit characters. As UTF-8 characters may be up to 32 bits the word rune is used.
}
```

switch i {
case 0: // empty case body
case 1:
f() // f is not called when i == 0!
}

```

And with:
```

switch i {
case 0: fallthrough
case 1:
f() // f is called when i == 0!
}

```

With default you can specify an action when none of the other cases match.
```

switch i {
case 0:
case 1:
f()
default:
g() // called when i is not 0 or 1
}

```

Cases can be presented in comma-separated lists.
```

func shouldEscape(c byte) bool {
switch c {
case ' ', '?', '\&', '=', '\#', '+': }\leftarrow, as "or"
return true
}
return false
}

```

Here's a comparison routine for byte arrays that uses two switch statements:
// Compare returns an integer comparing the two byte arrays
// lexicographically.
// The result will be 0 if \(a==b,-1\) if \(a<b\), and +1 if \(a>b\)
func Compare(a, b []byte) int \{
    for \(\mathrm{i}:=0\); \(\mathrm{i}<\operatorname{len}(\mathrm{a})\) \& \(\& \mathrm{i}<\operatorname{len}(\mathrm{b})\); i++ \(\{\)
        switch \{
        case \(a[i]>b[i]\) :
            return 1
        case \(a[i]<b[i]:\)
            return -1
        \}
    \}
    // String are equal except for possible tail
    switch \{
    case len(a) < len(b):
        return -1
    case len(a) > len(b):
        return 1
```

    }
    return 0 // Strings are equal
    }

```

\section*{Arrays, slices and maps}

Storing multiple values in a list can be done by utilizing arrays, or their more flexible cousin: slices. A dictionary or hash type is also available, it is called a map in Go.

\section*{Arrays}

An array is defined by: [ n ]<type>, where \(n\) is the length of the array. Assigning, or indexing an element in the array is done with square brackets:
```

var arr = [10]int
arr[0] = 42
arr[1] = 13
fmt.Printf("The first element is %d\n", arr[0])

```

Array types like var arr = [10] int have a fixed size. The size is part of the type. They can't grow, because then they would have a different type. Also arrays are values: Assigning one array to another copies all the elements. In particular, if you pass an array to a function, it will receive a copy of the array, not a pointer to it.

To declare an array you can use the following: var a [3]int, to initialize it to something else than zero, use a composite literal: a := [3]int\{1, 2, 3\} and this can be shortened to a := [...]int\{1, 2, 3\}, where Go counts the elements automatically. Note that all fields must be specified. So if you are using multidimensional arrays you have to do quite some typing
```

a := [2][2]int{ [2]int{1,2}, [2]int{3,4} }

```

Which is the same as:
```

a := [2][2]int{ [...]int{1,2}, [...]int{3,4} }

```

When declaring arrays you always have to type something in between the square brackets, either a number or three dots (. . .).

\section*{Slices}

A slice refers to an under laying array. What makes slices different from arrays is that a slice is a pointer to an array; slices are reference types, which means that if you assign one slice to another, both refer to the same underlying array. For instance, if a function takes a slice argument, changes it makes to the elements of the slice will be visible to the caller, analogous to passing a pointer to the underlying array. A slice is always coupled to an array that has a fixed size. For slices we define a capacity and a length.

Figure 2.1 depicts the following Go code. First we create an array of \(m+1\) elements of the type int: var array [m+1]int
Next, we create a slice from this array: slice := array[0:n+1]
And now we have:
- \(\operatorname{len}(s l i c e)==n, \operatorname{cap}(s l i c e)==m ;\)

A composite literal allows you to assign a value directly to an array, slice or map.

\section*{TODO}

Examples, with slice, slice[1:] slice[2:3] etc
- \(\operatorname{Len}(\) array \()==\boldsymbol{\operatorname { c a p }}(\) array \()==\mathrm{m}\).

Figure 2.1. Array versus slice


In the code listed in 2.6 we dare to do the impossible on line 8 and try to allocate something beyond the capacity (maximum length of the under laying array) and we are greeted with a runtime error.

Listing 2.6. Arrays and slices
\begin{tabular}{lll} 
package main & 1 \\
func main() \{ & & 3 \\
var array [100]int & // Create array, index from 0 to 99 & 4 \\
slice := array[0:99] & // Create slice, index from 0 to 98 & 5 \\
& & \\
slice[98] = 'a' & // OK & 7 \\
slice[99] = 'a' & // Error: "throw: index out of range" & 8 \\
\} &
\end{tabular}

Maps

Many other languages have a similar type built-in, Perl has hashes Python has its dictionaries and C++ also has maps (in lib) for instance. In Go we have the map type. A map can be thought of as an array indexed by strings (in its most simple form). In the following listing we define a map which converts from a string (month abbreviation) to an int - the number of days in that month. The generic way to define a map is with: map [<from type>]<to type>
```

monthdays := map[string]int{
"Jan": 31, "Feb": 28, "Mar": 31,
"Apr": 30, "May": 31, "Jun": 30,
"Jul": 31, "Aug": 31, "Sep": 30,
"Oct": 31, "Nov": 30, "Dec": 31, \leftarrow The comma here is required
}

```

For indexing (searching) in the map, we use square brackets, for example suppose we want to print the number of days in December: fmt. Printf("\%d\n", monthdays["Dec"])
If you are looping over an array, slice, string, or map a range clause help you again, which returns the key and corresponding value with each invocation.
```

year := 0
for _, days := range monthdays { // key is not used
year += days
}
fmt.Printf("Numbers of days in a year: %d\n", year)

```

Adding elements to the map would be done as:
```

monthday["Undecim"] = 30 // Add a month
monthday["Feb"] = 29 // Overwrite entry - for leap years

```

To test for existence, you would use the following[19]:
```

var value int
var present bool
value, present = monthday["Jan"] // Does it exist? If so present
has the value
// Or better and more Go like
v, ok := monthday["Jan"] // Hence, the "comma ok" form

```

And finally you can remove elements from the map:
monthday["Mar"] = 0, false // Deletes "Mar", always rains anyway
Which looks a bit like a reverse "comma ok" form.

\section*{Exercises}

Q2. (1) For-loop
1. Create a simple loop with the for construct. Make it loop 10 times and print out the loop counter with the fmt package.
2. Put the body of the loop in a separate function.
3. Rewrite the loop from 1. to use goto. The keyword for may not be used.

Q3. (1) FizzBuzz
1. Solve this problem, called the Fizz-Buzz [24] problem:

Write a program that prints the numbers from 1 to 100. But for multiples of three print "Fizz" instead of the number and for the multiples of five print "Buzz". For numbers which are multiples of both three and five print "FizzBuzz".

Q4. (1) Strings
1. Create a Go program that prints the following (up to 100 characters):

A
AA
AAA
AAAA
AAAAA
AAAAAA
AAAAAAA
2. Create a program that counts the numbers of characters \({ }^{e}\) in this string: asSASA ddd dsjkdsjs dk Make it also output the number of bytes in that string. Hint. Check out the utf8 package.
3. Extend the program from the previous question to replace the three runes at position 4 with 'abc'.
4. Write a Go program that reverses a string, so "foobar" is printed as "raboof".

\footnotetext{
\({ }^{e}\) In the UTF-8 world characters are sometimes called runes. Mostly, when people talk about characters, they
} mean 8 bit characters. As UTF-8 characters may be up to 32 bits the word rune is used.

\section*{Answers}

A2. (1) For-loop
1. There are a multitude of possibilities, one of the solutions could be:

Listing 2.7. Simple for-loop
package main
import "fmt"
func main() \{
for i := 0; i < 10; i++ \{
fmt.Printf("\%d\n", i)
\}
\}

Lets compile this and look at the output.
\% 6 g for.go \(\& \& 6 l-o\) for for. 8
\% ./for
0
1
\(\cdot\)

9
2. Next we put the body of the loop - the fmt. Printf - in a separate function.

Listing 2.8. Loop calls function
package main
import "fmt"
func main() \{
for i := 0; i < 10; i++ \{
show(i)
\}
\}
func show(j int) \{
fmt.Printf("\%d\n", j)
\}
The presented program should be self explanatory. Note however the " \(j\) int" instead of the more usual "int \(j\) " in the function definition.

A3. (1) FizzBuzz
1. A possible solution to this (relative) simple problem is the following program.

\section*{Listing 2.9. Fizz-Buzz}
```

package main
import "fmt"
func main() {
const (
FIZZ = 3
BUZZ = 5
)
for i := 1; i < 100; i++ {
p := false
if i%FIZZ == 0 {
fmt.Printf("Fizz")
p = true
}
if i%BUZZ == 0 {
fmt.Printf("Buzz")
p = true
}
if !p {
fmt.Printf("%v", i)
}
fmt.Println()
}
}

```

A4. (1) Strings
1. This program is a solution:

Listing 2.10. Strings
```

package main
import "fmt"
func main() {
str := "A"
for i := 0; i < 100; i++ {
fmt.Printf("%s\n", str)
str = str + "A"
}
}

```
2. To answer this question we need some help of the utf8 package. First we check the documentation with godoc utf8 | less. When we read the documentation we notice func RuneCount (p []byte)int. Secondly we can convert string to a byte slice with
```

str := "hello"
b := []byte(str) \leftarrowConversion, see page 50

```

Putting this together leads to the following program.

Listing 2.11. Runes in strings
package main
```

import (
"fmt"
"utf8"
)
func main() {
str := "dsjkdshdjsdh....js"
fmt.Printf("String %s\nLenght: %d, Runes: %d\n", str,
len([]byte(str)), utf8.RuneCount([]byte(str)))
}

```
3. Reversing a string can be done as follows:

Listing 2.12. Reverse a string
package main
func main() \{
s := "foobar"
a := []byte(s)
// Reverse a
for \(i, j:=0\), len(a)-1; \(i<j\) i, \(j=i+1, j-1\{\)
a[i], a[j] = a[j], a[i]
\}
println(string(a))
\}

\section*{3 Functions}

I'm always delighted by the light touch and stillness of early programming languages. Not much text; a lot gets done. Old programs read like quiet conversations between a well-spoken research worker and a well- studied mechanical colleague, not as a debate with a compiler. Who'd have guessed sophistication bought such noise?

RICHARD P. GABRIEL

Functions are the basic building blocks in Go programs; all interesting stuff happens in them. A function is declared as follows:

Listing 3.1. A function declaration

(1) The keyword func is used to declare a function;
(2) A function can be defined to work on a specific type, a more common name for such a function is method. This part is called a receiver and it is optional. See chapter 6;
(3) funcname is the name of your function;
(4) The variable q of type int is the input parameter. The parameters are passed pass-by-value meaning they are copied;

5 The variables \(r\) and \(s\) are the named return parameters for this function. Note that functions in Go can have multiple return values. See section "Multiple return values" on page 22 for more information. If you want the return parameters not to be named you only give the types: (int, int). If you have only one value to return you may omit the parentheses. If your function is a subroutine and does not have anything to return you may omit this entirely;

6 This is the function's body, note that return is a statement so the braces around the parameter(s) are optional.

Here are a two examples, the first is a function without a return value, the second is a simple function that returns its input.
```

func subroutine(in int) {
return
}

```
```

func identity(in int) int {
return in
}

```

Functions can be declared in any order you wish, the compiler scans the entire file before execution. So function prototyping is a thing of the past in Go.

Go disallows nested functions. You can however work around this by using anonymous functions, see section "Functions as values" on page 26 in this chapter.

\section*{Scope}

TODO
types in function, a nono. Maybe in package chapter.

Variables declared outside any functions are global in Go, those defined in functions are local to those functions. If names overlap - a local variable is declared with the same name as a global one - the local variable hides the global one when the current function is executed.

Listing 3.2. Local scope
Listing 3.3. Global scope


In listing 3.2 we introduce a local variable a in the function \(q()\). This local a is only visible in q() . That is why the code will print: 656. In listing 3.3 no new variables are introduced, there is only a global a. Assigning a new value to will be globally visible. This code will print: 655

In the following example we call \(g()\) from \(f()\) :

Listing 3.4. Scope when calling functions from functions

\section*{package main}
var a int
func main() \{
\(a=5\)
```

        println(a)
        f()
    }
func f() {
a := 6
println(a)
g()
}
func g() {
println(a)
}

```

The printout will be: 565. A local variable is only valid when we are executing the function in which it is defined.

\section*{Multiple return values}

One of Go's unusual features is that functions and methods can return multiple values (Python can do this too). This can be used to improve on a couple of clumsy idioms in C programs: in-band error returns (such as -1 for EOF) and modifying an argument.

In C , a write error is signaled by a negative count with the error code secreted away in a volatile location. In Go, Write returns a count and an error: "Yes, you wrote some bytes but not all of them because you filled the device". The signature of *File. Write in package os is:
func (file *File) Write(b []byte) (n int, err Error)
and as the documentation says, it returns the number of bytes written and a non-nil Error when \(n!=\operatorname{len}(b)\). This is a common style in Go.

A similar approach obviates the need to pass a pointer to a return value to simulate a reference parameter. Here's a simple-minded function to grab a number from a position in a byte array, returning the number and the next position.
```

func nextInt(b []byte, i int) (int, int) {
x := 0
// Naively assume everything is a number
for ; i < len(b); i++ {
x = x*10 + int(b[i])-'0'
}
return x, i
}

```

You could use it to scan the numbers in an input array a like this:
```

a := []byte{'1', '2', '3', '4'}
var x int
for i := 0; i < len(a); { < No i++
x, i = nextInt(a, i)
println(x)
}

```

Without having tuples as a native type, multiple return values is the next best thing to have. You can return precisely what you want without overloading the domain space with special values to signal errors.

\section*{Named result parameters}

The return or result parameters of a Go function can be given names and used as regular variables, just like the incoming parameters. When named, they are initialized to the zero values for their types when the function begins; if the function executes a return statement with no arguments, the current values of the result parameters are used as the returned values. Using this features enables you (again) to do more with less code \({ }^{\boldsymbol{a}}\).

The names are not mandatory but they can make code shorter and clearer: they are documentation. If we name the results of nextInt it becomes obvious which returned int is which.
```

func nextInt(b []byte, pos int) (value, nextPos int) { /* ... */ }

```

Because named results are initialized and tied to an unadorned return, they can simplify as well as clarify. Here's a version of io. ReadFull that uses them well:
```

func ReadFull(r Reader, buf []byte) (n int, err os.Error) {
for len(buf) > 0 \&\& err == nil {
var nr int
nr, err = r.Read(buf)
n += nr
buf = buf[nr:len(buf)]
}
return
}

```

In the following example we declare a simple function which calculates the factorial value of a value \(x\).

Some text in this section comes from [11].
```

func Factorial(x int) int { \leftarrow func Factorial(x int) (int) is also OK
if x == 0 {
return 1
} else {
return x * Factorial(x - 1)
}
}

```

So you could also write factorial as:
```

func Factorial(x int) (result int) {
if x == 0 {
result = 1
} else {
result = x * Factorial(x - 1)
}
return
}

```

\footnotetext{
\({ }^{a}\) This is a motto of Go; "Do more with less code".
}

When we use named result values, the code is shorter and easier to read. You can also write a function with multiple return values:
```

func fib(n) (val int, pos int) {
if n == 0 {
val = 1
pos = 0
} else if n == 1 {
val = 1
pos = 1
} else {
v1, _:= fib(n-1)
v2,_ := fib(n-2)
val = v1 + v2
pos = n
}
return
}

```

\section*{Deferred code}

Suppose you have a function in which you open a file and perform various writes and reads on it. In such a function there are often spots where you want to return early. If you do that, you will need to close the file descriptor you are working on. This often leads to the following code

\section*{Listing 3.5. Without defer}
```

func ReadWrite() bool {
file.Open("file")
// Do you thing
if failureX {
file.Close()
return false
}
if failureY {
file.Close()
return false
}
file.Close()
return true
}

```

Here a lot of code is repeated. To overcome this - and do more in less code - Go has the defer statement. After defer you specify a function which is called just before a return from the function is executed.

The code above could be rewritten as follows. This makes the function more readable, shorter and puts the Close right next to the Open.
```

func ReadWrite() bool {
file.Open("file")
defer file.Close() \leftarrow file.Close() is the function
// Do you thing
if failureX {
return false }\leftarrow\mathrm{ Close() is now done automatically
}
if failureY {
return false }\leftarrow\mathrm{ And here too
}
return true
}

```

You can put multiple functions on the "deferred list", like this example from [2]:
```

for i := 0; i < 5; i++ {
defer fmt.Printf("%d ", i)
}

```

Deferred functions are executed in LIFO order, so the above code prints: 43210 .
With defer you can even change return values, provided that you are using named result parameters and a function literal, i.e:

Listing 3.7. Function literal
defer func() \{
/* ... */
\(\}() \quad \leftarrow()\) is needed here
In that (unnamed) function you can access any named return parameter:
Listing 3.8. Access return values within defer
```

func f() (ret int) { \leftarrow ret is initialized with zero
defer func() {
ret++ }\leftarrow\mathrm{ Increment ret with 1
}()
return 0 }\leftarrow1\mathrm{ not 0 will be returned!
}

```

\section*{Variadic parameters}

Functions that take variadic parameters are functions that have a variable number of parameters. To do this, you first need to declare your function to take variadic arguments:

TODO
Rob Pike: Closures are just local functions

TODO
Extend this section a bit

Listing 3.9. Variadac function declaration
```

func myfunc(arg ... int) {}

```

The arg ... int instructs Go to see this as a function that takes a variable number of arguments. Note that these arguments all have the type int. Inside your function's body the variable arg is a slice of ints:
```

for _, n := range arg {
fmt.Printf("And the number is: %d\n", n)
}

```

If you don't specify the type of the variadic argument it defaults to the empty interface interface\{\} (see "??" in chapter 5).

\section*{Functions as values}

As with almost everything in Go, functions are also just values. They can be assigned to variables as follows:

Listing 3.10. Anonymous function
package main
```

func main() {
a := func() { // Define a nameless function and assign to a
println("Hello")
} // No () here
a() // Call the function
}

```

If we use fmt. Printf("\%T\n", a) to print the type of a, it prints func().
Functions-as-values may also be used in other places, like in maps. Here we convert from integers to functions:

Listing 3.11. Functions as values in maps
```

var xs = map[int]func() int{
func() int { return 10 },
func() int { return 20 },
func() int { return 30 },}\leqslant\mathrm{ Mandatory,
/* ... */
}

```

Or you can write a function that takes a function as its parameter, for example a Map function that works on int slices. This is left as an exercise for the reader, see exercise Q10 on page 27.

\section*{Panic and recovering}

\section*{Exercises}

Q5. (3) Integer ordering
1. Write a function that returns it parameters in the right, numerical (ascending) order:
\(\mathrm{f}(7,2) \rightarrow 2,7\)
\(\mathrm{f}(2,7) \rightarrow 2,7\)

Q6. (4) Scope
1. What is wrong with the following program?
package main
```

func main() { 5
for i := 0; i < 10; i++ { 6
fmt.Printf("%v\n", i) 7
} 8
fmt.Printf("%v\n", i) 9
}
10

```

Q7. (5) Stack
1. Create a simple stack which can hold a fixed amount of ints. It does not have to grow beyond this limit. Define both a push - put something on the stack - and a pop - retrieve something fro the stack - function. The stack should be a LIFO (last in, first out) stack.

Figure 3.1. A simple LIFO stack

2. Bonus. Write a String method which converts the stack to a string representation. This way you can print the stack using: fmt. Printf("My stack \%v\n", stack) The stack in the figure could be represented as: [0:m] [1:l] [2:k]
Q8. (5) Var args
1. Write a function that takes a variable numbers of ints and prints each integer on a separate line
Q9. (5) Fibonacci
BETTER
1. The Fibonacci sequence starts as follows: \(1,1,2,3,5,8,13, \ldots\) Or in mathematical terms: \(x_{1}=1 ; x_{2}=1 ; x_{n}=x_{n-1}+x_{n-2} \quad \forall n>2\).
Write a function that takes an int value and gives that many terms of the Fibonacci sequence.
Q10. (4) Map function \(A \operatorname{map}()\)-function is a function that takes a function and a list. The function is applied to each member in the list and a new list containing these calculated values is returned. Thus:
\(\operatorname{map}\left(f(),\left(a_{1}, a_{2}, \ldots, a_{n-1}, a_{n}\right)\right)=\left(f\left(a_{1}\right), f\left(a_{2}\right), \ldots, f\left(a_{n-1}\right), f\left(a_{n}\right)\right)\)
1. Write a simple map()-function in Go. It is sufficient for this function only to work for ints.
2. Expand your code to also work on a list of strings.

\section*{Answers}

A5. (3) Integer ordering
1.

A6. (4) Scope
1. The program does not even compile, because i on line 9 is not defined: i is only defined within the for-loop. To fix this the function main() should become:
```

func main() {
var i int
for i = 0; i < 10; i++ {
fmt.Printf("%v\n", i)
}
fmt.Printf("%v\n", i)
}

```

Now i is defined outside the for-loop and still visible after wards. This code will print the numbers 0 through 10.

A7. (5) Stack
1. First we define a new type that represents a stack; we need an array (to hold the keys) and an index, which points to the last element. Our small stack can only hold 10 elements.
```

type stack struct { }\leftarrow\mathrm{ stack is not exported
i int
data [10]int
}

```

Next we need the push and pop functions to actually use the thing. First we show the wrong solution! In Go data passed to functions is passed-by-value meaning a copy is created and given to the function. The first stab for the function push could be:
```

func (s stack) push(k int) { \leftarrow Works on copy of argument
if s.i+1 > 9 {
return
}
s.data[s.i] = k
s.i++
}

```

The function works on the s which is of the type stack. To use this we just call s.push (50), to push the integer 50 on the stack. But the push function gets a copy of \(s\), so it is not working the real thing. Nothing gets pushed to our stack this way, for example the following code:
```

var s stack \leftarrow make s a simple stack variable

```
s.push(25)
fmt.Printf("stack \%v\n", s);
s.push(14)
```

fmt.Printf("stack %v\n", s);
prints:
stack [0:0]
stack [0:0]

```

To solve this we need to give the function push a pointer to the stack. This means we need to change push from
func (s stack) push(k int) \(\rightarrow\) func (s *stack)push(k int)
We should now use new() (see "Allocation with new" in chapter 5) to create a pointer to a newly allocated stack, so line 1 from the example above needs to be \(s\) := new(stack)
And our two functions become:

Listing 3.12. The push and pop functions
```

func (s *stack) push(k int) {
s.data[s.i] = k
s.i++
}
func (s *stack) pop() int {
s.i-
return s.data[s.i]
}
Which we then use as follows

```

Listing 3.13. Stack usage
```

func main() {
var s stack
s.push(25)
s.push(14)
fmt.Printf("stack %v\n", s)
}

```
2. While this was a bonus question, having the ability to print the stack was very valuable when writing the code for this exercise. According to the Go documentation fmt. Printf("\%v") can print any value (\%v) that satisfies the Stringer interface. For this to work we only need to define a String() function for our type:

Listing 3.14. stack.String()
```

func (s stack) String() string {
var str string
for i := 0; i <= s.i; i++ {
str = str + "[" +
strconv.Itoa(i) + ":" + strconv.Itoa(s.data[i
]) + "]"

```
    \}
```

    return str
    }

```

A8. (5) Var args
1. For this we need the ...-syntax to signal we define a function that takes an arbitrary number of arguments.

Listing 3.15. A function with variable number of arguments
```

package main
import "fmt"
func main() {
printthem(1, 4, 5, 7, 4)
printthem(1, 2, 4)
}
func printthem(numbers ... int) { }\leftarrow\mathrm{ numbers is now a slice of ints
for _, d := range numbers {
fmt.Printf("%d\n", d)
}
}

```

A9. (5) Fibonacci
1. The following program calculates the Fibonacci numbers.

Listing 3.16. Fibonacci function in Go
```

package main
import "fmt"
func fibonacci(value int) []int {
x := make([]int, value)
x[0], x[1] = 1, 1 2
for n := 2; n < value; n++ {
x[n] = x[n-1] + x[n-2] 3
}
return x 4
}
func main() {
for _, term := range fibonacci(10) { (5)
fmt.Printf("%v ", term)
}
}

```
(1) We create an array to hold the integers up to the value given in the function call;
(2) Starting point of the Fibonicai calculation;
(3) \(x_{n}=x_{n-1}+x_{n-2}\);
(4) Return the entire array;

5 Using the range keyword we "walk" the numbers returned by the fibonacci funcion. Here up to 10. And we print them.

A10. (4) Map function

Listing 3.17. A Map function
1. func Map(f func(int) int, l []int) []int \{
j := make([]int, len(l))
for \(k, v:=\) range \(l\) \{ \(j[k]=f(v)\)
\}
return \(j\)
\}
func main() \{
m := []int\{1, 3, 4\}
\(\mathrm{f}:=\) func(i int) int \(\{\)
return i * i
\}
fmt.Printf("\%v", (Map(f, m)))
\}
2. Answer to question but now with strings

\section*{4 \\ Packages}

> Answer to whether there is a bitwise negation
> operator.

KEN THOMPSON

Packages are a collection of functions and data, they are like the Perl packages[8]. You declare a package with the package keyword. Note that (at least one of) the filename(s) should match the package name. Go packages may consist of multiple files, but they share the package <name> line. Lets define a package even in the file even.go.

Listing 4.1. A small package
```

package even }\leftarrow\mathrm{ Start our own namespace
func Even(i int) bool { < Exported function
return i % 2 == 0
}
func odd(i int) bool { }\leftarrow\mathrm{ Private function
return i % 2 == 1
}

```

Names that start with a capital letter are exported and may be used outside your package, more on that later. We can now use the package as follows in our own program myeven. go:

\section*{Listing 4.2. Use of the even package}
package main
import (
"./even"
"fmt"
)
func main() \{
i := 5 fmt.Printf("Is \%d even? \%v\n", i, even.Even(i))
\}
(1) Import the following packages;
(2) The local package even;
(3) The official fmt package;
(4) Use the function from the even package. The syntax for accessing a function from a package is <package>. Function.

Now we just need to compile and link, first the package, then myeven. go and then link it:
\begin{tabular}{lc}
\(\% 6 \mathrm{~g}\) even.go \(\quad\) \# package \\
\(\% 6 \mathrm{~g}\) myeven.go & \# our program \\
\(\% 6 \mathrm{l}\)-o myeven myeven. 6
\end{tabular}

And test it:
\% ./myeven
Is 5 even? false

In Go, a function from a package is exported (visible outside the package, i.e. public) when the first letter of the function name is a capital, hence the function name Even. If we change our myeven. go on line 7 to using to unexported function even. odd: fmt.Printf("Is \%d even? \%V\n", i, even.odd(i))

We get an error when compiling, because we are trying to use a private function:
```

myeven.go:7: cannot refer to unexported name even.odd
myeven.go:7: undefined: even.odd

```

To summarize:
- Public functions have a name starting with a capital letter;
- Private function have a name starting with a lowercase letter.

This convention also holds true for other names (new types, global variables) you define in a package.

\section*{Building a package}

The create a package that other people can use (by just using import "even") we first need to create a directory to put the package files in.
```

% mkdir even
% cp even.go even/

```

Next we can use the following Makefile, which is adapted for our even package.

Listing 4.3. Makefile for a package
```


# Copyright 2009 The Go Authors. All rights reserved.

```
1
\# Use of this source code is governed by a BSD-style ..... 2
\# license that can be found in the LICENSE file. ..... 3
include \$(GOROOT)/src/Make.inc ..... 5
TARG=even ..... 7
GOFILES=\} ..... 8
even.go\} ..... 9
include \$(GOROOT)/src/Make.pkg ..... 11

Note that on line 7 we define our even package and on the lines 8 and 9 we enter the files which make up the package. Also note that this is not the same Makefile setup as we used in section "Using a Makefile" in chapter 2. The last line with the include statement is different.

If we now issue gomake, a file named "_go_. 6 ", a directory named "_obj/" and a file inside " \({ }^{\circ} \mathrm{obj} /\) " called "even.a" is created. The file even. a is a static library which holds the compiled Go code. With gomake install the package (well, only the even.a) is installed in the official package directory:
```

% make install
cp _obj/even.a \$GOR00T/pkg/linux_amd64/even.a

```

After installing we can change our import from ". /even" to "even".
But what if you do not want to install packages in the official Go directory, or maybe you do not have the permission to do so? When using \(6 / 8 \mathrm{~g}\) you can use the \(I\)-flag to specify an alternate package directory. With this flag you can leave your import statement as-is (import "even") and continue development in your own directory. So the following commands should build (and link) our myeven program with our package.
```

% 6g -I even/_obj myeven.go \# building
% 6l -L even/_obj myeven.6 \# linking

```

\section*{Identifiers}

Names are as important in Go as in any other language. In some cases they even have semantic effect: for instance, the visibility of a name outside a package is determined by whether its first character is upper case. It's therefore worth spending a little time talking about naming conventions in Go programs.

The convention that is used was to leave well-known legacy not-quite-words alone rather than try to figure out where the capital letters go. Atoi, Getwd, Chmod. Camelcasing works best when you have whole words to work with: ReadFile, NewWriter, MakeSlice.

\section*{Package names}

When a package is imported (with import), the package name becomes an accessor for the contents. After
import "bytes"
the importing package can talk about bytes. Buffer. It's helpful if everyone using the package can use the same name to refer to its contents, which implies that the package name should be good: short, concise, evocative. By convention, packages are given lower case, single-word names; there should be no need for underscores or mixedCaps. Err on the side of brevity, since everyone using your package will be typing that name. And don't worry about collisions a priori. The package name is only the default name for imports. With the above import you can use bytes.Buffer. With
```

import bar "bytes"

```
it becomes bar. Buffer. So it does need not be unique across all source code, and in the rare case of a collision the importing package can choose a different name to use locally.

In any case, confusion is rare because the file name in the import determines just which package is being used.

Another convention is that the package name is the base name of its source directory; the package in src/pkg/container/vector is imported as container/vector but has name vector, not container_vector and not containerVector.

The importer of a package will use the name to refer to its contents, so exported names in the package can use that fact to avoid stutter. For instance, the buffered reader type in the bufio package is called Reader, not BufReader, because users see it as bufio.Reader, which is a clear, concise name. Moreover, because imported entities are always addressed with their package name, bufio. Reader does not conflict with io.Reader. Similarly, the function to make new instances of ring. Ring-which is the definition of a constructor in Go-would normally be called NewRing, but since Ring is the only type exported by the package, and since the package is called ring, it's called just New. Clients of the package see that as ring. New. Use the package structure to help you choose good names.

Another short example is once.Do; once.Do(setup) reads well and would not be improved by writing once. DoOrWaitUntilDone(setup). Long names don't automatically make things more readable. If the name represents something intricate or subtle, it's usually better to write a helpful doc comment than to attempt to put all the information into the name.

Finally, the convention in Go is to use MixedCaps or mixedCaps rather than underscores to write multiword names.

\section*{Initialization}

Every source file in a package can define an init() function. This function is called after the variables in the package have gotten their value. The init() function can be used to setup state before the execution begins.

\section*{Documenting packages}

Every package should have a package comment, a block comment preceding the package clause. For multi-file packages, the package comment only needs to be present in one file, and any one will do. The package comment should introduce the package and provide information relevant to the package as a whole. It will appear first on the godoc page and should set up the detailed documentation that follows. An example from the official regexp package:
```

/*
The regexp package implements a simple library for
regular expressions.
The syntax of the regular expressions accepted is:
regexp:
concatenation '|' concatenation
*/
package regexp

```

Each defined (and exported) function should has a small line of text documenting the behavior of the function, from the fmt package:
```

// Printf formats according to a format specifier and writes to standard output
func Printf(format string, a ...interface) (n int, errno os.Error) {

```

\section*{Testing packages}

In Go it is customary to write (unit) tests for your package. Writing tests involves the testing package and the program gotest. Both have excellent documentation. When you include tests with your package keep in mind that has to be build using the standard Makefile (see section "Building a package").

The testing itself is carried out with gotest. The gotest program run all the test functions. Without any defined tests for our even package a gomake test, yields:
```

% gomake test
no test files found (* test.go
make: *** [test] Error 2

```

Let us fix this by defining a test in a test file. Test files reside in the package directory and are named *_test.go. Those test files are just like other Go program, but gotest will only execute the test functions. Each test function has the same signature and its name should start with Test:
```

func TestXxx(t *testing.T) \leftarrow Test<Capital>restOftheName

```

When writing test you will need to tell gotest that a test has failed or was successful. A successful test function just returns. When the test fails you can signal this with the following functions [4]. These are the most important ones (see godoc testing for more):
func ( t *T) Fail()
Fail marks the test function as having failed but continues execution.
```

func (t *T) FailNow()

```

FailNow marks the test function as having failed and stops its execution. Execution will continue at the next test. So any other test in this file are skipped too
```

func (t *T) Log(args ...interface{})

```

Log formats its arguments using default formatting, analogous to Print(), and records the text in the error log.
```

func (t *T) Fatal(args ...interface{})

```

Fatal is equivalent to \(\log ()\) followed by FailNow().
Putting all this together we can write our test. First we pick a name: even_test.go. Then we add the following contents:

Listing 4.4. Test file for even package
package even
```

import "testing"3

```
```

func TestEven(t *testing.T) {5

```
        if true != Even(2) \{ 6
            t. Log("2 should be even!") 7
```

            t.Fail()8
    ```
\} ..... 9
\} ..... 10

Note that we use package even on line 1, the test fall in the same namespace as the package we are testing. This not only convenient, but also allows tests of unexported function and structures. We then import the testing package and on line 5 we define the only test function in this file. The displayed Go code should not hold any surprises: we check if the Even function works OK. Now, the moment we've been waiting for, executing the test:
```

% gomake test
6g -o _gotest_. 6 even.go even_test.go
rm -f _test/even.a
gopack grc _test/even.a _gotest_. }
PASS

```

Our test ran and reported PASS. Success! To show how a failed test look we redefine our test function:
```

// Entering the twilight zone
func TestEven(t *testing.T) {
if false != Even(2) {
t.Log("2 should be odd!")
t.Fail()
}
}

```

We now get:
--- FAIL: even.TestEven
2 should be odd!
FAIL
make: *** [test] Error 1

And you can act accordingly (by fixing the test for instance).
Writing new packages should go hand in hand with writing (some) documentation and test functions. It will make your code better and it shows that your really put in the effort.

\section*{Useful packages}

The standard Go repository includes a huge number of packages and it is even possible to install more along side your current Go installation. We cannot comment on each package, but the following a worth a mention: \({ }^{\boldsymbol{a}}\)
fmt
Package fmt implements formatted I/O with functions analogous to C's printf and scanf. The format verbs are derived from C's but are simpler. Some verbs (\%sequences) that can be used:

\footnotetext{
\({ }^{\boldsymbol{a}}\) The descriptions are copied from the packages' godoc
}
\(\% v\)
The value in a default format. when printing structs, the plus flag ( \(\%+\mathrm{v}\) ) adds field names;
\%\#v
a Go-syntax representation of the value.
\(\% T\)
a Go-syntax representation of the type of the value;
io
This package provides basic interfaces to I/O primitives. Its primary job is to wrap existing implementations of such primitives, such as those in package os, into shared public interfaces that abstract the functionality, plus some other related primitives.
bufio
This package implements buffered I/O. It wraps an io. Reader or io.Writer object, creating another object (Reader or Writer) that also implements the interface but provides buffering and some help for textual I/O.
sort
The sort package provides primitives for sorting arrays and user-defined collections.
strconv
The strconv package implements conversions to and from string representations of basic data types.
os
The os package provides a platform-independent interface to operating system functionality. The design is Unix-like.
flag
The flag package implements command-line flag parsing.
json
The json package implements encoding and decoding of JSON objects as defined in RFC 4627.

\section*{template}

Data-driven templates for generating textual output such as HTML.
Templates are executed by applying them to a data structure. Annotations in the template refer to elements of the data structure (typically a field of a struct or a key in a map) to control execution and derive values to be displayed. The template walks the structure as it executes and the "cursor" @ represents the value at the current location in the structure.
http
The http package implements parsing of HTTP requests, replies, and URLs and provides an extensible HTTP server and a basic HTTP client.
unsafe
The unsafe package contains operations that step around the type safety of Go programs.

\section*{reflect}

The reflect package implements run-time reflection, allowing a program to manipulate objects with arbitrary types. The typical use is to take a value with static type interface\{\} and extract its dynamic type information by calling Typeof, which returns an object with interface type Type. That contains a pointer to a struct of type *StructType, *IntType, etc. representing the details of the underlying type. A type switch or type assertion can reveal which (see chapter 6, section "Introspection").
exec
The exec package runs external commands.

\section*{Exercises}

Q11. (2) Stack as package See the 7 exercise. In this exercise we want to create a separate package for that code.
1. Create a proper package for your stack implementation, Push, Pop and the Stack type need to be exported.
2. Which official Go package could also be used for a (FIFO) stack?

Q12. (7) Calculator
1. Create a reverse polish calculator. Use your stack package.
2. Bonus. Rewrite your calculator to use the the package you found for question 2

\section*{Answers}

A11. (2) Stack as package
1. There are a few details that should be changed to make a proper package for our stack. First, the exported function should begin with a capital letter and so should Stack. So the full package (including the String() function becomes

\section*{Listing 4.5. Stack in a package}
```

package stack
import (
"strconv"
)
type Stack struct {
i int
data [10]int
}
func (s *Stack) Push(k int) {
s.data[s.i] = k
s.i++
}
func (s *Stack) Pop() (ret int) {
s.i--
ret = s.data[s.i]
}
func (s *Stack) String() string {
var str string
for i := 0; i < s.i; i++ {
str = str + "[" + strconv.Itoa(i) + ":"
+ strconv.Itoa(s.data[i]) + "]"
}
return str
}

```

A12. (7) Calculator
1. This is one answer

\section*{package main}
import (
"bufio"
```

    "os"
    "strconv"
    "fmt"
    )
var reader *bufio.Reader = bufio.NewReader(os.Stdin)
var st = new(Stack)
type Stack struct {
i int
data [10]int
}
func (s *Stack) push(k int) {
if s.i+1 > 9 {
return
}
s.data[s.i] = k
s.i++
}
func (s *Stack) pop() (ret int) {
s.i--
if s.i < 0 {
s.i = 0
return 0
}
ret = s.data[s.i]
return ret
}
func (s *Stack) String() string {
var str string
for i := 0; i < s.i; i++ {
str = str + "[" +
strconv.Itoa(i) + ":" + strconv.Itoa(s.data[i
]) + "]"
}
return str
}
func main() {
for {
s, err := reader.ReadString('\n')
var token string
if err != nil {
return
}
for _, c := range s {
switch {

```
```

            case c >= '0' && c <= '9':
                            token = token + string(c)
    case c == ' ':
r, _ := strconv.Atoi(token)
st.push(r)
token = ""
case c == '+':
fmt.Printf("%d\n", st.pop()+st.pop())
case c == '*':
fmt.Printf("%d\n", st.pop()*st.pop())
case c == '-''
p := st.pop()
q := st.pop()
fmt.Printf("%d\n", q-p)
case c == 'q':
return
default:
//error
}
}
}
}

```
2. The container/vector package would be a could candidate. It even comes with Push and Pop functions predefined. The changes needed to our program are minimal to say the least, the following unified diff shows the differences:
```

-- calc.go 2010-05-16 10:19:13.886855818 +0200
+++ calcvec.go 2010-05-16 10:13:35.000000000 +0200
@@ -5,11 +5,11 @@
"os"
"strconv"
"fmt"
"./stack"

+ "container/vector"
)
var reader *bufio.Reader = bufio.NewReader(os.Stdin)
-var st = new(Stack)
+var st = new(vector.IntVector)
func main() {
for {

```

Only two lines need to be changed. Nice.

\section*{5 \\ Beyond the basics}

Go has pointers but not pointer arithmetic. You cannot use a pointer variable to walk through the bytes of a string.

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Go has pointers. There is however no pointer arithmetic, so they act more like references than pointers that you may know from C. Pointers are are useful. Remember that when you call a function in Go, the variables are pass-by-value. So, for efficiency and the possibility to modify a passed value in functions we have pointers.

Just like in C you declare a pointer by prefixing the type with an '*': var \(p\) *int. Now \(p\) is a pointer to an integer value. All newly declared variables are assigned their zero value and pointers are no difference. A newly declared, or just a pointer that points to nothing has a nil-value. In other languages this is often called a NULL pointer in Go it is just nil. To make a pointer point to something you can use the address-of operator ( \(\&\) ), which we do on line 5 :

Listing 5.1. Make use of a pointer
```

var $p$ *int 1
fmt.Printf("\%v", p) $\leftarrow$ Prints nil 2
var i int $\quad \leftarrow$ Declare integer variable $i \quad 4$
$\mathrm{p}=\& \mathrm{i} \quad \leftarrow$ Make $p$ point to $i \quad 5$
fmt.Printf("\%v", p) $\leftarrow$ Prints something like 0x7ff96b81c000a 7

```

More general: *X is a pointer to an \(X\); [3] \(X\) is an array of three \(X\) s. The types are therefore really easy to read just read out the names of the type modifiers: [] declares an array slice; '*' declares a pointer; [size] declares an array. So []*[3]*x is an array slice of pointers to arrays of three pointers to Xs (also see figure 5.1).

Figure 5.1. Pointers and types, the values \(v\) all have type \(x\)
\begin{tabular}{ll} 
*X & \(\rightarrow v\) \\
{\([3] X\)} & {\(\left[v_{\theta}\right]\left[v_{1}\right]\left[v_{2}\right]\)} \\
{[]\(X\)} & {\(\left[v_{\theta}\right] \cdots\left[v_{2}\right]\)} \\
{[]\(*[3] * X\)} & {\(\left[n_{\theta}\right] \cdots\left[n_{2}\right]\)} \\
& {\(\left[k_{\theta}\right] \rightarrow v_{\theta}\)} \\
& {\(\left[k_{1}\right] \rightarrow v_{1}\)} \\
& {\(\left[k_{2}\right] \rightarrow v_{2}\)}
\end{tabular}

Dereferencing a pointer is done by prefixing the pointer variable with '*':

\section*{Listing 5.2. Dereferencing a pointer}
\begin{tabular}{ll}
\(\mathrm{p}=\& \mathrm{i}\) & \(\leftarrow\) Take the address of \(i\) \\
\(* \mathrm{p}=8\) & \(\leftarrow\) Change the value of \(i\) \\
fmt.Printf("\%v\n", *p) & \(\leftarrow\) Prints 8 \\
fmt.Printf("\%v\n", i) & \(\leftarrow\) Idem
\end{tabular}

As said, there is no pointer arithmetic, so if you write: \({ }^{*} p++\), it is interpreted as \((* p)++\) : first deference and then increment the value.

\section*{Allocation}

Go has garbage collection, meaning that you don't have to worry about memory allocation and deallocation. Of course almost every language since 1980 has this, but it is nice to see garbage collection in a C-like language.

Go has two allocation primitives, new and make. They do different things and apply to different types, which can be confusing, but the rules are simple. The following sections show how to handle allocation in Go and hopefully clarifies the somewhat artificial distinction between new and make.

\section*{Allocation with new}

The built-in function new is essentially the same as its namesakes in other languages: new (T) allocates zeroed storage for a new item of type \(\mathbf{T}\) and returns its address, a value of type *T. In Go terminology, it returns a pointer to a newly allocated zero value of type \(\mathbf{T}\). This is important to remember:
| new returns pointers.

This means a user of the data structure can create one with new and get right to work. For example, the documentation for bytes.Buffer states that "the zero value for Buffer is an empty buffer ready to use." Similarly, sync. Mutex does not have an explicit constructor or Init method. Instead, the zero value for a sync.Mutex is defined to be an unlocked mutex.

The zero-value-is-useful property works transitively. Consider this type declaration.
```

type SyncedBuffer struct {
lock sync.Mutex
buffer bytes.Buffer
}

```

Values of type SyncedBuffer are also ready to use immediately upon allocation or just declaration. In this snippet, both p and v will work correctly without further arrangement.
```

p := new(SyncedBuffer) // type *SyncedBuffer
var v SyncedBuffer // type SyncedBuffer

```

\section*{Allocation with make}

Back to allocation. The built-in function make( \(T\), args) serves a purpose different from new ( \(T\) ). It creates slices, maps, and channels only, and it returns an initialized (not zero) value of type \(T\), not \(* T\). The reason for the distinction is that these three types are, under the covers, references to data structures that must be initialized before use. A slice, for
example, is a three-item descriptor containing a pointer to the data (inside an array), the length, and the capacity; until those items are initialized, the slice is nil. For slices, maps, and channels, make initializes the internal data structure and prepares the value for use.

I make returns initialized (non zero) values.
For instance, make([]int, 10, 100) allocates an array of 100 ints and then creates a slice structure with length 10 and a capacity of 100 pointing at the first 10 elements of the array. In contrast, new ( [ ]int) returns a pointer to a newly allocated, zeroed slice structure, that is, a pointer to a nil slice value.

These examples illustrate the difference between new () and make().
```

var p *[]int = new([]int) // allocates slice structure; *p == nil
// rarely useful
var v []int = make([]int, 100) // v refers to a new array of 100 ints
// Unnecessarily complex:
var p *[]int = new([]int)
*p = make([]int, 100, 100)
// Idiomatic:
v := make([]int, 100)

```

Remember that make() applies only to maps, slices and channels and does not return a pointer. To obtain an explicit pointer allocate with new().

Constructors and composite literals

Sometimes the zero value isn't good enough and an initializing constructor is necessary, as in this example taken from the package os.
```

func NewFile(fd int, name string) *File {
if fd < 0 {
return nil
}
f := new(File)
f.fd = fd
f.name = name
f.dirinfo = nil
f.nepipe = 0
return f
}

```

There's a lot of boiler plate in there. We can simplify it using a composite literal, which is an expression that creates a new instance each time it is evaluated.
```

func NewFile(fd int, name string) *File {
if fd < 0 {
return nil
}
f := File{fd, name, nil, 0} \leftarrowCreate a new File
return \&f }\leftarrow\mathrm{ Return the address of }
}

```

Note that it's perfectly OK to return the address of a local variable; the storage associated with the variable survives after the function returns. In fact, taking the address of a composite literal allocates a fresh instance each time it is evaluated, so we can combine these last two lines.
return \&File\{fd, name, nil, 0\}
The fields of a composite literal are laid out in order and must all be present. However, by labeling the elements explicitly as field:value pairs, the initializers can appear in any order, with the missing ones left as their respective zero values. Thus we could say
return \&File\{fd: fd, name: name\}
As a limiting case, if a composite literal contains no fields at all, it creates a zero value for the type. The expressions new (File) and \&File\{\} are equivalent.

Composite literals can also be created for arrays, slices, and maps, with the field labels being indices or map keys as appropriate. In these examples, the initializations work regardless of the values of Enone, Eio, and Einval, as long as they are distinct.
```

ar := [...]string {Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}
sl := []string {Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}
ma := map[int]string{Enone: "no error", Eio: "Eio", Einval: "invalid
argument"}

```

\section*{Defining your own}

Of course Go allows you to define new types, it does this with the type keyword:
```

type foo int

```

Creates a new type foo which acts like an int. Creating more sophisticated types is done with the struct keyword. An example would be when we want record somebody's name (string) and age (int) in a single structure and make it a new type:

Listing 5.3. Structures
```

package main
import "fmt"
type NameAge struct {
name string }\quad\leftarrow\mathrm{ Not exported
age int \leftarrowNot exported, Age would be exported
}
func main() {
a := new(NameAge)
a.name = "Pete"
a.age = 42
fmt.Printf("%v\n", a)
}

```

Apropos, the output of fmt.Printf("\%v\n", a) is

Pete, 42

That is nice! Go knows how to print your structure. If you only want to print one, or a few, fields of the structure you'll need to use .<field name>. For example to only print the name:
fmt.Printf("\%s", a.name) \(\leftarrow \%\) formats a string
TODO
Struct with exported fields

Methods

If you create functions that works on your newly defined type, you can take two routes:
1. Create a function that takes the type as an argument.
func doSomething(in1 *NameAge, in2 int) \{ /* ... */ \}

This is (you might have guessed) a function call.
2. Create a function that works on the type (see receiver in listing 3.1):
func (in1 *NameAge) doSomething(in2 int) \{ /* ... */ \}

This is a method call, which can be used as:
var n *NameAge
n.doSomething(2)

But keep the following in mind, this is quoted from [3]:

If \(\boldsymbol{x}\) is addressable and \(\& x\) 's method set contains \(m, x . m()\) is shorthand for ( \(\& x\) ). \(m()\).

In the above case this means that the following is not an error:
```

var n NameAge }\leftarrowNot a pointer
n.doSomething(2)

```

Here Go will search the method list for \(n\) of type NameAge, come up empty and will then also search the method list for the type *NameAge and will translate this call to (\&n).doSomething (2).

\section*{Conversions}

Sometimes you want to convert a type to another type. In C this is known as casting a value to another type. This is also possible in Go, but there are some rules. For starters, converting from one value to another is done by functions and not all conversions are allowed.

Table 5.1. Valid conversions
\begin{tabular}{|c|c|c|c|c|c|}
\hline From & xb []byte & xi []int & s string & f float & i int \\
\hline \multicolumn{6}{|l|}{To} \\
\hline []byte & \(\times\) & & []byte(s) & & \\
\hline []int & & \(\times\) & []int(s) & & \\
\hline string & string(xb) & string(xi) & \(\times\) & & \\
\hline float & & & & \(\times\) & float(i) \\
\hline int & & & & int (f) & \(\times\) \\
\hline
\end{tabular}
- From a string to a slice of bytes or ints.
```

mystring := "hello this is string"
byteslice := []byte(mystring)

```

Converts to a byte slice, each byte contains the integer value of the corresponding byte in the string. Note that as strings in Go are encoded in UTF-8 some characters in the string may end up in \(1,2,3\) or 4 bytes.
```

intslice := []int(mystring)

```

Converts to an int slice, each int contains a Unicode code point. Every character from the string is corresponds to one integer.
- From a slice of bytes or ints to a string.
```

b := []byte{'h','e','l','l','o'} \leftarrowComposite literal
s := string(b)
i := []int{257,1024,65}
r := string(i)

```

For numeric values the following conversion are defined:
- Convert to a integer with a specific (bit) length: uint8(int);
- From floating point to an integer value: int(float). This discards the fraction part from the floating point value;
- The other way around: float(int);

User defined types

TODO
We create two types here Foo and Bar, where Bar is an alias for Foo:
```

type foo struct { int }
type bar foo

```

Then we:
```

var b bar = bar{1} }\leftarrow\mathrm{ Declare b to be a bar
var f foo = b }\quad\leftarrow\mathrm{ Assign b to f

```

Make the text flow more nicely into this paragraph.

TODO
CHECK THIS.

Which fails on the last line with:
cannot use b (type bar) as type foo in assignment
This can be fixed with a conversion:
var f foo \(=\mathrm{foo}(\mathrm{b})\)

\section*{Exercises}

Q13. (6) Map function with interfaces Use the answer from exercise Q10, but now make it generic using interfaces.
Q14. (6) Pointers
1. Suppose we have defined the following structure:
```

type Person struct {
name string
age int
}

```

What is the difference between the following two lines?
```

var p1 Person

```
p2 := new(Person)
2. What is the difference between the following two allocations?
```

func Set(t *T) {
x = t
}
and
func Set(t T) {
x= \&t
}

```

Q15. (5) Pointers and reflection
1. One of the last paragraphs in section "Introspection and reflection" on page 61, has the following words:

The code on the right works OK and sets the member Name to "Albert Einstein". Of course this only works when you call Set() with a pointer argument.

Why is this the case?
Q16. (6) Linked List
1. Make use of the package container/list to create a (double) linked list. Push the values 1,2 and 4 to the list and then print it.
2. Create your own linked list implementation. And perform the same actions is in question 1
Q17. (6) Cat
1. Write a program which mimics Unix program cat. For those who don't know this program, the following invocation displays the contents of the file blah:
2. Make it support the \(n\) flag, where each line is numbered.

Q18. (8) Method calls
1. Suppose we have the following program:
```

package main
import "container/vector"
func main() {
k1 := vector.IntVector{}
k2 := \&vector.IntVector{}
k3 := new(vector.IntVector)
k1.Push(2)
k2.Push(3)
k3.Push(4)

```
\}

What are the types of k1, k2 and k3?
2. Now, this program compiles and runs OK. All the Push operations work even though the variables are of a different type. The documentation for Push says:
func ( \(p^{*}\) IntVector) Push( \(x\) int) Push appends \(x\) to the end of the vector.
So the receiver has to be of type *IntVector, why does the code above work then?

\section*{Answers}

A13. (6) Map function with interfaces
Listing 5.4. A generic map function in Go
```

package main
import "fmt"
/* define the empty interface as a type */
type e interface{}
func mult2(f e) e {
switch f.(type) {
case int:
return f.(int) * 2
case string:
return f.(string) + f.(string) + f.(string) + f.(string)
}
return f
}
func Map(n []e, f func(e) e) []e {
m := make([]e, len(n))
for k, v := range n {
m[k] = f(v)
}
return m
}
func main() {
m := []e{1, 2, 3, 4}
s := []e{"a", "b", "c", "d"}
mf := Map(m, mult2)
sf := Map(s, mult2)
fmt.Printf("%V\n", mf)
fmt.Printf("%v\n", sf)
}

```

A14. (6) Pointers
1. In first line: var p1 Person allocates a Person-value to p1. The type of p1 is Person. The second line: p2 := new(Person) allocates memory and assigns a pointer to p2. The type of p 2 is *Person.
2. In the second function, \(x\) points to a new (heap-allocated) variable \(t\) which contains a copy of whatever the actual argument value is.
In the first function, \(x\) points to the same thing that \(t\) does, which is the same thing that the actual argument points to.
So in the second function, we have an "extra" variable containing a copy of the interesting value.

A15. (5) Pointers and reflection
1. When called with a non-pointer argument the variable is a copy (call-by-value). So you are doing the reflection voodoo on a copy. And thus you are not changing the original value, but only this copy.

A16. (6) Linked List
1.
2.

A17. (6) Cat
1. A solution might be:

Listing 5.5. A cat program
```

package main
import (
"os"
"fmt"
"bufio"
"flag"
)
var numberFlag = flag.Bool("n", false, "number each line")
func cat(r *bufio.Reader) {
i := 1
for {
buf, e := r.ReadBytes('\n')
if e == os.EOF {
break
}
if *numberFlag {
fmt.Fprintf(os.Stdout, "%5d %s", i, buf)
i++
} else {
fmt.Fprintf(os.Stdout, "%s", buf)
}
}
return
}
func main() {
flag.Parse()
if flag.NArg() == 0 {
cat(bufio.NewReader(os.Stdin))
}
for i := 0; i < flag.NArg(); i++ {
f, e := os.Open(flag.Arg(i), os.0_RDONLY, 0)

```
```

        if e != nil {
        fmt.Fprintf(os.Stderr, "%s: error reading from
        %s: %s\n",
                        os.Args[0], flag.Arg(i), e.String())
    continue
    }
cat(bufio.NewReader(f))
}
}

```

A18. (8) Method calls
1. The type of \(k 1\) is vector. IntVector. Why? We use a composite literal (the \{\}), so we get a value of that type back. The variable k2 is of *vector. IntVector, because we take the address ( \(\&\) ) of the composite literal. And finally k3 has also the type *vector. IntVector, because new returns a pointer to the type.
2. The answer is given in [3] in the section "Calls", where among other things it says:

A method call \(x . m()\) is valid if the method set of (the type of) \(\times\) contains \(m\) and the argument list can be assigned to the parameter list of \(m\). If \(x\) is addressable and \&x's method set contains m, x.m() is shorthand for ( \(\delta x\) ) . \(m\) ().

In other words because k1 is addressable and *vector.IntVector does have the Push method, the call k1.Push(2) is translated by Go into (\&k1).Push(2) which makes the type system happy again (and you too - now you know this). \({ }^{\boldsymbol{a}}\)

\footnotetext{
\({ }^{a}\) Also see section "Methods" in this chapter.
}

\section*{6 Inetifaes}

I have this phobia about having my body penetrated surgically. You know what I mean?
eXistenZ
TED PIKUL

The following text is from [23]. Written by Ian Lance Taylor one of the authors of Go.

In Go, the word interface is overloaded to mean several different things. Every type has an interface, which is the set of methods defined for that type. This bit of code defines a struct type \(\mathbf{S}\) with one field, and defines two methods for \(\mathbf{s}\).

Listing 6.1. Defining a struct and methods on it
```

type S struct { i int }
func (p *S) Get() int { return p.i }
func (p *S) Put(v int) { p.i = v }

```

You can also define an interface type, which is simply a set of methods. This defines an interface I with two methods:
```

type I interface {
Get() int
Put(int)
}

```
| An interface type is a set of methods.
\(\mathbf{S}\) is a valid implementation for \(\mathbf{I}\), because it defines the two methods which \(\mathbf{I}\) requires. Note that this is true even though there is no explicit declaration that \(\mathbf{S}\) implements I. A Go program can use this fact via yet another meaning of interface, which is an interface!value:
```

func f(p I) { fmt.Println(p.Get()); p.Put(1) }

```

Here the variable p holds a value of interface type. Because \(\mathbf{S}\) implements \(\mathbf{I}\), we can call \(f\) passing in a pointer to a value of type \(\mathbf{s}\) :
```

var s S; f(\&s)

```

The reason we need to take the address of \(\mathbf{s}\), rather than a value of type \(\mathbf{s}\), is because we defined the methods on s to operate on pointers, see the code above in listing 6.1. This is not a requirement - we could have defined the methods to take values - but then the Put method would not work as expected.

The fact that you do not need to declare whether or not a type implements an interface means that Go implements a form of duck typing[27]. This is not pure duck typing, because when possible the Go compiler will statically check whether the type implements the interface. However, Go does have a purely dynamic aspect, in that you can convert from one interface type to another. In the general case, that conversion is checked at runtime. If the conversion is invalid - if the type of the value stored in the existing interface value does not satisfy the interface to which it is being converted - the program will fail with a runtime error.

Interfaces in Go are similar to ideas in several other programming languages: pure abstract virtual base classes in C++, typeclasses in Haskell or duck typing in Python. However there is no other language which combines interface values, static type checking, dynamic runtime conversion, and no requirement for explicitly declaring that a type satisfies an interface. The result in Go is powerful, flexible, efficient, and easy to write.

\section*{Empty interface}

For example, since every type satisfies the empty interface: interface \{\}. We can create a generic function which has an empty interface as its argument:

Listing 6.2. A function with a empty interface argument
```

func g(any interface{}) int {
return any.(I).Get()
}

```

The return any.(I).Get() is the tricky bit in this function. The value any has type interface, meaning no guarantee of any methods at all: it could contain any type. The .(I) is a type switch which converts any to an interface of type I. If we have that type we can invoke the Get() function. So if we create a new variable of the type *s, we can just call \(g()\), because *S also implements the empty interface.
```

s = new(S)
fmt.Println(g(s));

```

The call to g will work fine and will print 0 . If we however invoke g() with a value that does not implement I we have a problem:

Listing 6.3. Failing to implement an interface
```

i := 5 // make i a "lousy" int

```
fmt. Println(g(i))

This compiles OK, but when we run this we get slammed with:
panic: interface conversion: int is not main.I: missing method Get Which is completely true, the built-in type int does not have a Get () function.

\section*{Methods}

Methods are functions that have an receiver (see chapter 3). You can define methods on any type (except the built-ins like int) can have methods. You can make an integer type with its own methods. For example:
```

type Foo int
func (self Foo) Emit() {
fmt.Printf("%v", self)
}
type Emitter interface {
Emit()
}

```

Doing this on built-in (are types defined in other package) types yields:

Listing 6.4. Failure extending built-in types
```

func (i int) Emit() {
fmt.Printf("%d", i)
}
cannot define new methods
on non-local type int

```

Listing 6.5. Failure extending non-local types
```

func (a *net.AddrError) Emit() {
fmt.Printf("%v", a)
}
cannot define new methods
on non-local type net.AddrError

```

Another thing about the receiver; it can not be a defined for interface types, doing so results in a invalid receiver type ... compiler error. The authoritative word from the language spec [3]:

The receiver type must be of the form \(\boldsymbol{T}\) or \(\boldsymbol{T}^{\boldsymbol{T}}\) where T is a type name. T is called the receiver base type or just base type. The base type must not be a pointer or interface type and must be declared in the same package as the method.

\section*{Interface names}

By convention, one-method interfaces are named by the method name plus the -er suffix: Reader, Writer, Formatter etc.

There are a number of such names and it's productive to honor them and the function names they capture. Read, Write, Close, Flush, String and so on have canonical signatures and meanings. To avoid confusion, don't give your method one of those names unless it has the same signature and meaning. Conversely, if your type implements a method with the same meaning as a method on a well-known type, give it the same name and signature;

\section*{Introspection}

In a program, you can discover the dynamic type of an interface variable by using a switch. Such a type switch uses the syntax of a type assertion with the keyword type inside the parentheses. If the switch declares a variable in the expression, the variable will have the corresponding type in each clause.

Listing 6.6. Dynamically find out the type
```

package main
type PersonAge struct { 1
name string
age int
}
type PersonShoe struct { 2
name string
shoesize int
}
func main() {
p1 := new(PersonAge)
p2 := new(PersonShoe)

```
```

    WhichOne(p1)
    WhichOne(p2)
    }
func WhichOne(x interface{}) {
switch t := x.(type) {4
case *PersonAge:
O
println("Age person")
case *PersonShoe:
println("Shoe person")
}
}

```
(1) First we define two structures as a new type, PersonAge;

2 And PersonShoe;
(3) This function must accept both types as valid input, so we use the empty Interface, which every type implements;
(4) The type switch: (type);
(5) When allocated with new it's a pointer. So we check for *PersonAge. If WhichOne() was called with a non pointer value, we should check for PersonAge.

The following is another example of performing a type switch, but this time checking for more (built-in) types:

Listing 6.7. A more generic type switch
```

switch t := interfaceValue.(type) { }\leftarrow\mathrm{ The type switch

```
case bool:
    fmt.Printf("boolean \%t\n", t)
case int:
    fmt.Printf("integer \%d\n", t)
case *bool:
    fmt.Printf("pointer to boolean \%t n ", *t)
case *int:
    fmt.Printf("pointer to integer \%d n ", *t)
default:
    fmt. Printf("unexpected type \%T", t) // \%T prints type
\}

Introspection and reflection

In the following example we want to look at the "tag" (here named "namestr") defined in the type definition of Person. To do this we need the reflect package (there is no other way in Go). Keep in mind that looking at a tag means going back the type definition. So we use the reflectpackage to figure out the type of the variable and then access the tag.

\section*{Listing 6.8. Introspection using reflection}
```

type Person struct {
name string "namestr"
age int
}
p1 := new(Person) \leftarrow new returns a pointer to Person
ShowTag(p1) \leftarrowShowTag() is now called with this pointer
func ShowTag(i interface{}) {
switch t := reflect.NewValue(i).(type) { < Type assert on reflect value
case *reflect.PtrValue: \leftarrow Hence the case for *reflect.PtrValue
tag := t.Elem()}.\mp@subsup{\underbrace}{}{\mathrm{ Type()}}\cdot(*\mp@subsup{\underbrace}{}{*reflect.StructType)}. F\mp@subsup{\underbrace}{}{\mathrm{ Field(0)}}\cdot\mp@subsup{\underbrace}{}{\mathrm{ Tag}

```

(1) We are dealing with a PtrValue and according to the documentation \({ }^{\boldsymbol{a}}\) :
func (v *PtrValue) Elem() Value
Elem returns the value that v points to. If v is a nil pointer, Elem returns a nil Value.
we can use Elem() to get the type the pointer points to. In this case *reflect. StructValue. We have also used reflect.NewValue(i) for the type assertion, so that we get back types in the *reflect namespace;

2 On a Value we can use the function Type() which returns reflect. Type;
(3) Again according to the documentation, we have:
...which returns an object with interface type Type. That contains a pointer to a struct of type *StructType, *IntType, etc. representing the details of the underlying type. A type switch or type assertion can reveal which.

So we can access your specific type as a member of this struct. Which we do with (*reflect.StructType);
(4) A StructType has a number of methods, one of which is Field ( \(n\) ) which returns the \(n^{t h}\) field of a structure. The type of this return is a StructField;

5 The struct StructField has a Tag member which returns the tag-name as a string. So on the \(0^{\text {th }}\) field we can unleash . Tag to access this name: Field(0).Tag. This finally gives us namestr.

To make the difference between looking a types and values more clear, that a look at the following code:

Listing 6.9. Reflection and the type and value
func show(i interface\{\}) \{
\(\boldsymbol{a}_{\text {godoc }}\) reflect
```

switch t := i.(type) {
case *Person:
r := reflect.NewValue(i) \leftarrow Enter the world of reflection
tag :=
r.(*reflect.PtrValue).Elem().Type().(*reflect.StructType).Field
(0).tag
nam := 2
r.(*reflect.PtrValue).Elem().(*reflect.StructValue).Field(0)।
newline.(*reflect.StringValue).Get()
}

```
\}
(1) Here we want to get the "tag", which means going for the type. Thus we need Elem().Type(). (*reflect.StructType) to get to it;
(2) Now we want to get access to the value of one of the members and we employ Elem(). (*reflect.StructValue) to get to it. Now we have arrived at the structure. Then we go the the first field Field(0), tell reflect is a *reflect.StringValue and invoke the Get () method on it.

Figure 6.1. Pealing away the layers using reflection. Going from \(a\) *Person via *reflect.PtrValue using the methods described in godoc reflect to get the actual string contained deep within.


Reflection works by pealing off layers once you have got your hands on a value in the reflection world.

Setting a value works similarly as getting a value, but only works on exported members. Again some code:

Listing 6.10. Reflect with private member
```

type Person struct {

```
    name string "namestr"
    age int
\}
    func Set(i interface\{\}) \{
    switch t := i.(type) \{
    case *Person:
        \(r:=\) reflect.NewValue(i)
        r.(*reflect.PtrValue).Elem().
        (*reflect.StructValue).
        FieldByName("name").
        (*reflect.StringValue).
        Set("Albert Einstein")
        \}
\}

Listing 6.11. Reflect with public member
```

type Person struct {

```
    Name string "namestr" \(\leftarrow\)
    age int
\}
func Set(i interface\{\}) \{
    switch \(\mathrm{t}:=\mathrm{i}\).(type) \{
    case *Person:
    \(r:=\) reflect.NewValue(i)
    r.(*reflect.PtrValue).Elem().
    (*reflect.StructValue).
    FieldByName("Name"). \(\leftarrow\)
    (*reflect.StringValue).
    Set("Albert Einstein")
    \}
\}

The code on the left compiles and runs, but when you run it, you are greeted with a stack trace and a runtime error:
panic: cannot set value obtained via unexported struct field
The code on the right works OK and sets the member Name to "Albert Einstein". Of course this only works when you call Set() with a pointer argument.

\section*{Exercises}

Q19. (6) Interfaces and compilation
1. The code in listing 6.3 on page 59 compiles OK - as stated in the text. But when you run it you'll get a runtime error, so something is wrong. Why does the code compile cleanly then?

\section*{Answers}

A19. (6) Interfaces and compilation
1. The code compiles because an integer type implements the empty interface and that is the check that happens at compile time.
A proper way to fix this, is to test if such an empty interface can be converted and if so, call the appropriate method. The Go code that defines the function g in listing 6.2 - repeated here:
func \(g(a n y\) interface\{ \(\}\) ) int \(\{\) return any.(I). Get() \}
Should be changed to become:
func \(g(a n y\) interface\{\}) int \{
if \(v\), ok := any.(I); ok \{ // Check if any can be converted
return v.Get() // If so invoke Get()
\}
return -1 // Just so we return anything
\}
If g() is called now there are no run-time errors anymore. The idiom used is called "comma ok" in Go.

\section*{7 \\ Concurrency}
- Parallelism is about performance;
- Concurrency is about program design.

Google IO 2010
ROBE PIKE

In this chapter we will show off Go's ability for concurrent programming using channels and goroutines. Goroutines are the central entity in Go’s ability for concurrency. But what is a goroutines? From [2]:

They're called goroutines because the existing terms - threads, coroutines, processes, and so on - convey inaccurate connotations. A goroutine has a simple model: it is a function executing in parallel with other goroutines in the same address space. It is lightweight, costing little more than the allocation of stack space. And the stacks start small, so they are cheap, and grow by allocating (and freeing) heap storage as required.

A goroutine is a normal function, except that you start it with the keyword go.
```

ready("Tee", 2) \leftarrowNormal function call
go ready("Tee", 2) \leftarrowReady() started as goroutine

```

The following idea for a program was taken from [20]. We run a function as two goroutines, the goroutines wait for an amount of time and them print something to the screen. On the lines 14 and 15 we start the goroutines. The main function waits long enough, so that both goroutines will have printed their text. Right now we wait for 5 seconds (time. Sleep() counts in ns) on line 17, but in fact we have no idea how long we should wait until all goroutines have exited.

Listing 7.1. Go routines in action
```

func ready(w string, sec int) {8

```
time.Sleep(int64(sec) * 1e9) ..... 9
fmt. Println(w, "is ready!") ..... 10
\} ..... 11
func main() \{ ..... 13
go ready("Tee", 2) ..... 14
go ready("Coffee", 1) ..... 15
fmt. Println("I'm waiting") ..... 16
time.Sleep(5 * 1e9) ..... 17
\} ..... 18

Listing 7.1 outputs:
I'm waiting \# right away
Coffee is ready! \# after 1 minute
Tee is ready! \# after 2 minutes

If we did not wait for the goroutines (i.e. remove line 17) the program would be terminated immediately and any running goroutines would die with it. To fix this we need some kind of mechanism which allows us to communicate with the goroutines. This mechanism is available to us in the form of channels. A channel can be compared to a two-way pipe in Unix shells: you can send to and receive values from it. Those values can only be of a specific type: the type of the channel. If we define a channel, we must also define the type of the values we can send on the channel. Note that we must use make to create a channel:
```

ci := make(chan int)
CS := make(chan string)
cf := make(chan interface{})

```

Makes ci a channel on which we can send and receive integers, makes cs a channel for strings and cf a channel for types that satisfy the empty interface. Sending on a channel and receiving from it, is done with the same operator: <-. Depending on the operands it figures out what to do:
```

ci<- 1 < Send the integer 1 to the channel ci
<-ci }\leftarrow\mathrm{ Receive an integer from the channel ci
i :=<-ci \leftarrow Receive from the channel ci and storing it in i

```

Lets put this to use.
Listing 7.2. Go routines and a channel
```

var c chan int 1 1
func ready(w string, sec int) { 3
time.Sleep(int64(sec) * le9) 4
fmt.Println(w, "is ready!") 5
c<- 1 2 6
} 7
func main() {i 9
c = make(chan int) 3 10
go ready("Tee", 2) 4 11
go ready("Coffee", 1) 12
fmt.Println("I'm waiting, but not too long") 13
<-c 5
<-c 6
15
} 16

```
(1) Declare c to be a variable that is a channel of ints. That is: this channel can move integers. Note that this variable is global so that the goroutines have access to it;
(2) Send the integer 1 on the channel c;
(3) Initialize c;
(4) Start the goroutines with the keyword go;
(5) Wait until we receive a value from the channel. Note that the value we receive is discarded;

TODO
Do this with range,
use select later

6 Two goroutines, two values to receive

There is still some remaining ugliness; we have to read twice from the channel (lines 14 and 15). This is OK in this case, but what if we don't know how many goroutines we started? This is where another Go built-in comes in: select. With select you can (among other things) listen for incoming data on a channel.

Using select in our program does not really make it shorter, because we run too few goroutines. We remove the lines 14 and 15 and replace them with the following:

\section*{Listing 7.3. Using select}
```

L: for { 14
select { 15
case <-c: 16
i++ 17
if i > 1 { 18
break L 19
} 20
} 21
} 22

```

Make it run in parallel
While our goroutines were running concurrent, they were not running in parallel. When you do not tell Go anything there can only be one goroutine running at a time. With runtime. \(\operatorname{GOMAXPROCS}(\mathrm{n})\) you can set the number of goroutines that can run in parallel. From the documentation:

GOMAXPROCS sets the maximum number of CPUs that can be executing simultaneously and returns the previous setting. If \(n<1\), it does not change the current setting. This call will go away when the scheduler improves.

If you do not want to change any source code you can also set an environment variable GOMAXPROCS to the desired value.

\section*{Exercises}

Q20. (4) Channels
1. Modify the program you created in exercise Q2 to use channels, in other words, the function called in the body should now be a goroutine and communication should happen via channels. You should not worry yourself on how the goroutine terminates.
2. There are a few annoying issues left if you resolve question 1 . One of the problems is that the goroutine isn’t neatly cleaned up when main.main() exits. And worse, due to a race condition between the exit of main.main() and main.shower() not all numbers are printed. It should print up until 9, but sometimes it prints only to 8 . Adding a second quit-channel you can remedy both issues. Do this. \({ }^{a}\)
Q21. (7) Fibonacci II

\footnotetext{
\({ }^{a}\) You will need the select statement.
}
1. This is the same exercise as the one given page 27 in exercise 9 . For completeness the complete question:

The Fibonacci sequence starts as follows: \(1,1,2,3,5,8,13, \ldots\) Or in mathematical terms: \(x_{1}=1 ; x_{2}=1 ; x_{n}=x_{n-1}+x_{n-2} \quad \forall n>2\).
Write a function that takes an int value and gives that many terms of the Fibonacci sequence.
But now the twist: You must use channels.

\section*{Answers}

A20. (4) Channels
1. A possible program is:

\section*{Listing 7.4. Channels in Go}
```

package main 1
import "fmt" 3
func main() { 5
ch := make(chan int) 6
go shower(ch) 7
for i := 0; i < 10; i++ { 8
ch <- i 9
} 10
} 11
func shower(c chan int) { 13
for { 14
j := <-c 15
fmt.Printf("%d\n", j) 16
} 17
} 18

```

We start of in the usual way, then at line 6 we create a new channel of ints. In the next line we fire off the function shower with the ch variable as it argument, so that we may communicate with it. Next we start our for-loop (lines 8-10) and in the loop we send (with <-) our number to the function (now a goroutine) shower. In the function shower we wait (as this blocks) until we receive a number (line 15). Any received number is printed (line 16) and then continue the endless loop started on line 14.
2. An answer is

Listing 7.5. Adding an extra quit channel
package main ..... 1
import "fmt" ..... 3
func main() \{ ..... 5
ch := make(chan int) ..... 6
quit := make(chan bool) ..... 7
go shower(ch, quit) ..... 8
for i := 0; i < 10; i++ \{ ..... 9
ch <- i ..... 10
\} ..... 11
quit <- false // or true, does not matter ..... 12
```

}15

```
```

func shower(c chan int, quit chan bool) {

```
```

func shower(c chan int, quit chan bool) {

```
for \{ ..... 16
select \{ ..... 17
case j := <-c: ..... 18
fmt.Printf("\%d\n", j) ..... 19
case <-quit: ..... 20
break ..... 21
\} ..... 22
\} ..... 23
\} ..... 24

On line 20 we read from the quit channel and we discard the value we read. We could have used q := <-quit, but then we would have used the variable only once - which is illegal in Go. Another trick you might have pulled out of your hat may be: _ = <-quit. This is valid in Go, but the Go idiom favors the one given on line 20.

A21. (7) Fibonacci II
1. The following program calculates the Fibonacci numbers using channels.

\section*{Listing 7.6. A Fibonacci function in Go}
```

package main
import "fmt"
func dup3(in <-chan int) (<-chan int, <-chan int, <-chan int) {
a, b, c := make(chan int, 2), make(chan int, 2), make(chan int
, 2)
go func() {
for {
x := <-in
a <- x
b <- x
c <- x
}
}()
return a, b, c
}
func fib() <-chan int {
x := make(chan int, 2)
a, b, out := dup3(x)
go func() {
x <- 0
x <- 1
<-a
for {
x<- <-a+<-b

```
```

                                    }
        }()
        return out
    }
func main() {
x := fib()
for i := 0; i < 10; i++ {
fmt.Println(<-x)
}
}

```
// See sdh33b.blogspot.com/2009/12/fibonacci-in-go.html

\section*{8 \\ Communication}

Good communication is as stimulating as black coffee, and just as hard to sleep after.

ANNE MORROW LINDBERGH

In this chapter we are going to look at the building blocks in Go for communicating with the outside world.

\section*{Files}

Reading from (and writing to) files is easy in Go. This program only uses the os package to read data from the file /etc/passwd.

\section*{Listing 8.1. Reading from a file (unbufferd)}
\begin{tabular}{|c|c|}
\hline package main & 1 \\
\hline import "os" & 3 \\
\hline func main() \{ & 5 \\
\hline buf := make([]byte, 1024) & 6 \\
\hline f, _ := os.Open("/etc/passwd", os.0_RDONLY, 0666) & 7 \\
\hline defer f.Close() & 8 \\
\hline for \(\{\) & 9 \\
\hline n, _ \(:=\mathrm{f}\).Read(buf) & 10 \\
\hline if \(\mathrm{n}==0\) \{ & 11 \\
\hline break & 12 \\
\hline \} & 13 \\
\hline os.Stdout.Write(buf[0:n]) & 14 \\
\hline \} & 15 \\
\hline \} & 16 \\
\hline
\end{tabular}

If you want to use buffered IO there is the bufio package:

Listing 8.2. Reading from a file (bufferd)
package main ..... 1
import ( ..... 3
"os" ..... 4
"bufio" ..... 5
) ..... 6
func main() \{ ..... 8
buf := make([]byte, 1024) ..... 9
f, _:= os.Open("/etc/passwd", os.0_RDONLY, 0666) ..... 10
defer f.Close() ..... 11
\(r\) := bufio.NewReader(f) ..... 12
```

    w := bufio.NewWriter(os.Stdout) 13
    defer w.Flush() 14
    for { 15
    n, _:= r.Read(buf) 16
    if n == 0 { 17
            break 18
            } 19
            w.Write(buf[0:n]) 20
            } 21
    }19

```
w.Write(buf[0:n]) ..... 20
\} ..... 21
\} ..... 22

On line 12 we create a bufio. Reader from f which is of type *File. NewReader expects an io. Reader, so you might think this will fail. But it doesn't. An io. Reader is defined as:
```

type Reader interface {
Read(p []byte) (n int, err os.Error)
}

```

So anything that has such a Read() function implements this interface. And from listing 8.1 (line 10) we can see that *File indeed does so.

\section*{Command line arguments}

Arguments from the command line are available inside your program via the string slice os. Args, provided you have imported the package os. The flag package has a more sophisticated interface, and also provided a way to parse flags.

TODO
example.

\section*{Executing commands}

The exec package has function to run external commands, and it the premier way to execute commands from within a Go program. We start commands with the Run function:
```

func Run(argv0 string, argv, envv []string, dir string, stdin, stdout,
stderr int) (p *Cmd, err os.Error)

```

Run starts the binary prog running with arguments argv and environment envv. It returns a pointer to a new Cmd representing the command or an error.

Lets execute ls \(-l\) :
import "exec"
cmd, err := exec.Run("/bin/ls", []string\{"ls", "-l"\}, nil, "", exec. DevNull, exec.DevNull, exec.DevNull)

In the os package we find the ForkExec function. This is another way (but more low level) to start executables. \({ }^{\boldsymbol{a}}\) The prototype for ForkExec is:
func ForkExec(argv0 string, argv []string, envv []string, dir string, fd []*File) (pid int, err Error)

\footnotetext{
\({ }^{\boldsymbol{a}}\) There is talk on the go-nuts mailing list about separating Fork and Exec.
}

With the following documentation:
ForkExec forks the current process and invokes Exec with the file, arguments, and environment specified by argv0, argv, and envv. It returns the process id of the forked process and an Error, if any. The fd array specifies the file descriptors to be set up in the new process: fd[0] will be Unix file descriptor 0 (standard input), fd[1] descriptor 1, and so on. A nil entry will cause the child to have no open file descriptor with that index. If dir is not empty, the child chdirs into the directory before execing the program.

Suppose we want to execute ls -l again:
```

import "os"
pid, err := os.ForkExec("/bin/ls", []string{"ls", "-l"}, nil, "", []*os.
File{ os.Stdin, os.Stdout, os.Stderr})
defer os.Wait(pid, os.WNOHANG) \leftarrow Otherwise you create a zombie

```

\section*{Exercises}

Q22. (8) Processes
1. Write a program that takes a list of all running processes and prints how many child processes each parent has spawned. The output should look like:
```

Pid 0 has 2 children: [1 2]
Pid 490 has 2 children: [1199 26524]
Pid 1824 has 1 child: [7293]

```
- For acquiring the process list, you'll need to capture the output of ps -e -opid, ppid, comm. This output looks like:
\begin{tabular}{rl} 
PID & PPID COMMAND \\
9024 & 9023 zsh \\
19560 & 9024 ps
\end{tabular}
- If a parent has one child you must print child, is there are more than one print children;
- The process list must be numerically sorted, so you start with pid 0 and work your way up.

Here is a Perl version to help you on your way (or to create complete and utter confusion).

Listing 8.3. Processes in Perl
```

\#!/usr/bin/perl -l
my (%child, \$pid, $parent);
my @ps=`ps -e -opid,ppid,comm'; # Capture the ouput from 'ps'
foreach (@ps[1..$\#ps]) { \# Discard the header line
(\$pid, $parent, undef) = split; # Split the line, discard 'comm'
    push @{$child{\$parent}}, \$pid; \# Save the child PIDs on a list
}

```
```


# Walk through the sorted PPIDs

foreach (sort { \$a <=> \$b } keys %child) {
print "Pid ", $_, " has ", @{$child{$_}}+0, " child", # Print them
        @{$child{$_}} == 1 ? ": " : "ren: ", "[@{$child{\$_}}]";
}

```

\section*{Answers}

A22. (8) Processes
1. There is lots of stuff to do here. We can divide our program up in the following sections:
1. Starting ps and capturing the output;
2. Parsing the output and saving the child PIDs for each PPID;
3. Sorting the PPID list;
4. Printing the sorted list to the screen

In the solution presented below, we've opted to use container/vector to hold the PIDs. This "list" grows automatically.
The function atoi (lines 19 through 22) is defined to get ride of the multiple return values of the original strconv. Atoi, so that it can be used inside function calls that only accept one argument, as we do on lines 45,47 and 50 .
A possible program is:
Listing 8.4. Processes in Go
```

package main 1
import (3

```
"os" ..... 4
"fmt" ..... 5
"sort" ..... 6
"bufio" ..... 7
"strings" ..... 8
"strconv" ..... 9
"container/vector" ..... 10
) ..... 11
const ( ..... 13
PID = iota ..... 14
PPID ..... 15
) ..... 16
func atoi(s string) (x int) \{ ..... 18
\(x, \quad=\) strconv.Atoi (s) ..... 19
return ..... 20
\} ..... 21
func main() \{ ..... 23
pr, pw, _ := os.Pipe() ..... 24
defer pr.Close() ..... 25
r := bufio.NewReader(pr) ..... 26
w := bufio.NewWriter(os.Stdout) ..... 27
defer w.Flush() ..... 28
pid, _ := os.ForkExec("/bin/ps", []string\{"ps", "-e", "-opid, ..... 29
```ppid,comm"\}, nil, "", []*os.File\{nil, pw, nil\})
```

defer os.Wait(pid, os.WNOHANG) ..... 30
pw.Close() ..... 31
child := make(map[int]*vector.IntVector) ..... 33
s, ok := r.ReadString('\n') // Discard the header line ..... 34
$s, o k=r . R e a d S t r i n g(' \backslash n ')$ ..... 35
for ok == nil \{ ..... 36
f := strings.Fields(s) ..... 37
if _, present := child[atoi(f[PPID])]; !present \{ ..... 38
v := new(vector.IntVector) ..... 39
child[atoi(f[PPID])] = v ..... 40
\} ..... 41
// Save the child PIDs on a vector ..... 42
child[atoi(f[PPID])].Push(atoi(f[PID])) ..... 43
s, ok = r.ReadString('\n') ..... 44
\} ..... 45
// Sort the PPIDs ..... 47
schild := make([]int, len(child)) ..... 48
i := 0 ..... 49
for $k$, $\quad$ := range child \{ ..... 50
schild[i] = k ..... 51
i++ ..... 52
\} ..... 53
sort.SortInts(schild) ..... 54
// Walk throught the sorted list ..... 55
for _, ppid := range schild \{ ..... 56
fmt. Printf("Pid \%d has \%d child", ppid, child[ppid]. ..... 57
Len())
if child[ppid].Len() == 1 \{ ..... 58
fmt.Printf(": \%v\n", []int(*child[ppid])) ..... 59
\} else \{ ..... 60
fmt.Printf("ren: \%v\n", []int(*child[ppid])) ..... 61
\} ..... 62
\} ..... 63

## Exercises

These exercises need to be put in "their" chapter. They are included here so the at least validate as $\mathbb{a r}_{\mathrm{E}} \mathrm{X}$.

## Exercises

Q23. (3) Minimum and maximum

1. Write a function that calculates the maximum value in an int slice ([]int).
2. Write a function that calculates the minimum value in a int slice ([]int).

Q24. (5) Bubble sort

1. Write a function that performs Bubble sort on slice of ints. From [25]:

It works by repeatedly stepping through the list to be sorted, comparing each pair of adjacent items and swapping them if they are in the wrong order. The pass through the list is repeated until no swaps are needed, which indicates that the list is sorted. The algorithm gets its name from the way smaller elements "bubble" to the top of the list.
Q25. (5) Word and letter count

1. Write a small program that reads text from standard input and performs the following actions:
2. Count the number of characters (including spaces);
3. Count the number of words;
4. Count the numbers of lines.

In other words implement wc (1) (check you local manual page), however you only have to read from standard input.

Q26. (4) Average

1. Write a function that calculates the average of a float slice.

Q27. (9) Number cruncher

- Pick six (6) random numbers from this list:

$$
1,2,3,4,5,6,7,8,9,10,25,50,75,100
$$

Numbers may be picked multiple times;

- Pick one (1) random number ( $i$ ) in the range: 1 . . . 1000;
- Tell how, by combining the first 6 numbers (or a subset thereof) with the operators ,,$+- *$ and /, you can make $i$;

An example. We have picked the numbers: $1,6,7,8,8$ and 75 . And $i$ is 977 . This can be done in many different ways, one way is:

$$
((((1 * 6) * 8)+75) * 8)-7=977
$$

or

$$
(8 *(75+(8 * 6)))-(7 / 1)=977
$$

1. Implement a number cruncher that works like that. Make it print the solution in a similar format (i.e. output should be infix with parenthesis) as used above.
2. Calculate all possible solutions and show them (or only show how many there are). In the example above there are 544 ways to do it.

## Q28. (4) Uniq

1. Write a Go program that mimics the function of the Unix uniq command. This program should work as follows, given a list with the following items:
'a' 'b' 'a' 'a' 'a' 'c' 'd' 'e' 'f' 'g'
it should print only those item which don't have the same successor:
'a' 'b' 'a' 'c' 'd' 'e' 'f'
Listing A. 4 is a Perl implementation of the algorithm.
Listing A.4. uniq(1) in Perl
```
#!/usr/bin/perl
my @a = qw/a b a a a c d e f g/;
print my $first = shift @a;
foreach (@a) {
    if ($first ne $_) { print; $first = $_; }
}
```


## Answers

A23. (3) Minimum and maximum

1. This the function for calculating a maximum:
```
func max(l []int) (max int) {
    max = l[0]
    for _, v := range l {
                if v > max {
                max = v
            }
        }
        return
}
```

2. This the function for calculating a minimum:
```
func min(l []int) (min int) {
    min = l[0]
    for _, v := range l {
                if v < min {
                min = v
            }
        }
        return
}
```

A24. (5) Bubble sort
1.

A25. (5) Word and letter count

1. The following program is an implementation of $\mathrm{wc}(1)$.

Listing A.1. wc(1) in Go

```
package main
import (
    "os"
    "fmt"
    "bufio"
    "strings"
)
func main() {
    var chars, words, lines int
    r := bufio.NewReader(os.Stdin)
    for {
                switch s, ok := r.ReadString('\n'); true {
```

```
    case ok != nil:
    fmt.Printf("%d %d %d\n", chars, words, lines);
    return
    default:
        chars += len(s)
        words += len(strings.Fields(s))
        lines++
}
        }
}
```

A26. (4) Average

1. The following function calculates the average.

Listing A.2. Average function in Go
func average(xs []float) (ave float) \{
sum := 0.0
switch len(xs) \{
case 0:
ave $=0$
default:
for _, v := range xs \{
sum += v
\}
ave = sum / float(len(xs))
\}
return
\}
func main() \{
\}

A27. (9) Number cruncher

Listing A.3. Number cruncher
1.

| package main |  |
| :---: | :---: |
| import |  |
|  | "fmt" |
|  | "strconv" |
|  | "container/vector" |
|  | "flag" |
| $)$ |  |
| const ( |  |
|  | $\underset{\text { ADD }}{=} 1000 * \text { iota }$ |
|  | SUB |
|  | MUL |
|  | DIV |
|  | MAXPOS $=11$ |
| ) |  |

```
var mop = map[int]string{
    ADD: "+",
    SUB:"-",
    MUL: "*",
    DIV: "/",
}
var (
    ok bool
    value int
)
type Stack struct {
    i int
    data [MAXPOS]int
}
func (s *Stack) Reset() {
    s.i = 0
}
func (s *Stack) Len() int {
    return s.i
}
func (s *Stack) Push(k int) {
    s.data[s.i] = k
    s.i++
}
func (s *Stack) Pop() int {
    s.i..
    return s.data[s.i]
}
var found int
var stack = new(Stack)
func main() {
    flag.Parse()
    list := []int{1, 6, 7, 8, 8, 75, ADD, SUB, MUL, DIV}
    // list := []int1, 6, 7, ADD, SUB, MUL, DIV
    magic, ok := strconv.Atoi(flag.Arg(0))
    if ok != nil {
            return
    }
    f := make([]int, MAXPOS)
    solve(f, list, 0, magic)
}
func solve(form, numberop []int, index, magic int) {
    var tmp int
    for i, v:= range numberop {
        if v == 0 {
            goto NEXT
        }
            if v < ADD {
            // it is a number, save it
            tmp = numberop[i]
            tmp = numberop[i]
            }
            form[index] = v
            value, ok = rpncalc(form[0 : index+1])
            if ok &&& value == magic {
            if v < ADD {
                numberop[i] = tmp // reset and go on
            }
            found++
            fmt.Printf("%s = %d #%d\n", rpnstr(form[0:index+1]), value, found)
            //goto NEXT
            }
            if index == MAXPOS-1 {
            if v < ADD {
                numberop[i] = tmp // reset and go on
            }
            }
            solve(form, numberop, index+1, magic)
            if v < ADD {
            numberop[i] = tmp // reset and go on
            }
```

```
    NEXT:
}
/ convert rpn to nice infix notation and string
// the r must be valid rpn form
func rpnstr(r []int) (ret string) {
    := new(vector.StringVector)
    or k, t := range r {
        switch t {
        case ADD, SUB, MUL, DIV
            a := s.Pop()
            := s.Pop()
            if k== Len(r)-1{
                s.Push(b + mop[t] + a)
                } else {
                    s.Push("(" + b + mop[t] + a + ")"
        default:
            s.Push(strconv.Itoa(t))
        }
    for
        := range *s {
        ret += v
    }
}
// return result from the rpn form
| if the expression is not valid, ok is false
func rpncalc(r []int) (int, bool) {
    stack.Reset()
        for_, t := range r
            witch t {
                case ADD, SUB, MUL, DIV:
```



```
                a := stack.Pop()
                b := stack.Pop()
                if t == ADD {
                    stack.Push(b + a)
                if t== SUB {
                    / disallow negative subresults
                    if b-a<0 {
                                    return 0, false
                    }
                    stack.Push(b - a)
                }
                f t == MUL {
                    stack.Push(b * a)
                }
                f t == DIV {
                    if a == 0 {
                            }
                            // disallow fractions
                    if b%a != 0 {
                                    return 0, false
                    }
                    stack.Push(b / a)
                }
                default:
                stack.Push(t)
                }
    }
    if stack.Len() == 1 { // there is only one.
        return stack.Pop(), true
    }
    return 0, false
}
```

2. When starting permrec we give 977 as the first argument:

| \% ./permrec 977 |  |
| :--- | :--- |
| $1+(((6+7) * 75)+(8 / 8))=977$ | $\# 1$ |
| $\ldots$ | $\ldots$ |
| $((75+(8 * 6)) * 8)-7=977$ | $\# 542$ |
| $(((75+(8 * 6)) * 8)-7) * 1=977$ | $\# 543$ |
| $(((75+(8 * 6)) * 8)-7) / 1=977$ | $\# 544$ |

A28. (4) Uniq

1. The following is a uniq implementation in Go.

Listing A.5. uniq(1) in Go

```
package main
import "fmt"
func main() {
    list := []string{"a", "b", "a", "a", "c", "d", "e", "f"}
        // fmt.Printf("first := list[0]
        fmt.Printf("%s ", first)
        for _, v := range list[1:] {
            if first != v {
                fmt.Printf("%s ", v)
                first = v
            }
        }
}
```


## B Colophon

This work was created with $\square_{\mathrm{E}} \mathrm{X}$. The main text is set in the Google Droid fonts. All typewriter text is typeset in DejaVu Mono.

## Contributors

The following people have helped to make this book what it is today. In no particular order:
Miek Gieben [miek@miek.nl](mailto:miek@miek.nl)
Main text, exercises and answers;
JC van Winkel
Proof reading, exercises and changes to the main text.
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## Further reading

All text used to create this book are listed in the bibliography, but here follows a more detailed list of the most interesting ones.

Detailed explanation about defer and the panic and recover mechanism in Go
$\rightarrow$ http://blog.golang.org/2010/08/defer-panic-and-recover.html.
array
capacity, 13
length, 13
multidimensional, 13
buffered, 74
built-in
cap, 7
close, 6
closed, 6
cmplx, 7
copy, 7
imag, 7
len, 7
make, 7, 47
new, 7, 47
panic, 7
panicln, 7
print, 7
println, 7
real, 7
channel, 67
channels, 1, 67
complex numbers, 7
composite literal, 13
deferred list, 25
duck typing, 58
function
as values, 26
call, 50
literal, 25
literals, 26
goroutine, 66
goroutines, 1
gotest, 38
interface, 58
set of methods, 58
type, 58
value, 58
keyword
break, 8, 10
continue, 10
default, 12
defer, 24
else, 8
fallthrough, 11
for, 9
go, 66
goto, 9
if, 8
import, 36
iota, 5
map, 14
add elements, 15
existence, 15
remove elements, 15
package, 34
range, 10, 14
on maps, 11,14
on slices, 11
return, 8
select, 68
struct, 49
switch, 11
type, 49
method, 20
method call, 50
MixedCaps, 37
named return parameters, 20
nil, 46
operator
address-of, 46
and, 6
bit clear, 6
channel, 67
increment, 47
or, 6
xor, 6
package
bufio, 37, 40, 74
bytes, 36
container/vector, 37
even, 34
exec, 41,75
flag, 40
fmt, 7, 39
http, 40
io, 40
json, 40
lib, 14
os, 40,75
reflect, 41, 61
ring, 37
sort, 40
strconv, 40
template, 40
unsafe, 40
parallel assignment, 4, 10
pass-by-value, 20
private, 35
public, 35
receiver, 20
reference types, 13
runes, 11, 16
scope
local, 21
slice
capacity, 13
length, 13
type switch, 59, 60
variables
, 4
assigning, 3
declaring, 3
underscore, 4

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[^0]:    $\boldsymbol{a}_{\text {http://golang.org/doc/ itself is served by a Go program called godoc. }}$

[^1]:    ${ }^{a}$ Yes, that sounds a lot like coroutines, but goroutines are slightly different as we will see in chapter 7.

